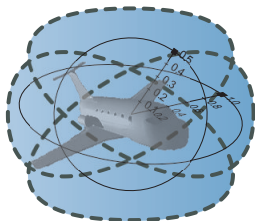


Logic & Proofs for Cyber-Physical Systems with KeYmaera X

André Platzer

Carnegie Mellon University





- 1 CPS are Multi-Dynamical Systems
 - Hybrid Systems / Games / Stochastic / Distributed Hybrid Systems
- 2 Differential Dynamic Logic
- 3 Axioms and Proofs for CPS
- 4 Differential Invariants for Differential Equations
 - Differential Invariants
 - Example: Elementary Differential Invariants
- 5 Applications
 - Ground Robot Navigation
 - Airborne Collision Avoidance System
 - KeYmaera X
- 6 Summary



Which control decisions are safe for aircraft collision avoidance?

Cyber-Physical Systems

CPSs combine cyber capabilities with physical capabilities to solve problems that neither part could solve alone.

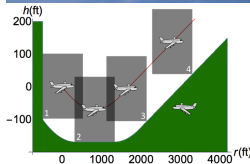
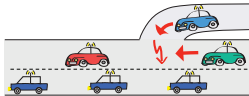
CPSs Promise Transformative Impact!

Prospects: Safe & Efficient

Driver assistance
Autonomous cars

Pilot decision support
Autopilots / UAVs

Train protection
Robots near humans



Prerequisite: CPSs need to be safe

How do we make sure CPSs make the world a better place?

Can you trust a computer to control physics?

Can you trust a computer to control physics?

- 1 Depends on how it has been programmed
- 2 And on what will happen if it malfunctions

Rationale

- 1 Safety guarantees require analytic foundations.
- 2 A common foundational core helps all application domains.
- 3 Foundations revolutionized digital computer science & our society.
- 4 Need even stronger foundations when software reaches out into our physical world.

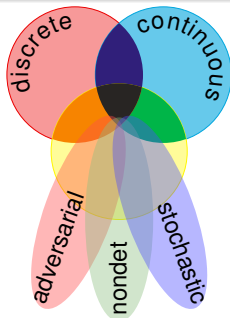
CPSs deserve proofs as safety evidence!



CPSs are Multi-Dynamical Systems

CPS Dynamics

CPS are characterized by multiple facets of dynamical systems.



CPS Compositions

CPS combines multiple simple dynamical effects.

Descriptive simplification

Tame Parts

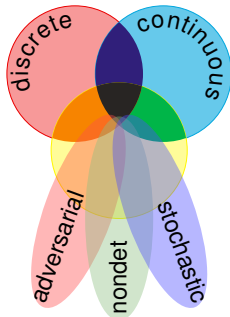
Exploiting compositionality tames CPS complexity.

Analytic simplification

CPSs are Multi-Dynamical Systems

hybrid systems

HS = discrete + ODE

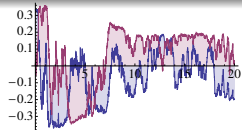


hybrid games

HG = HS + adversary

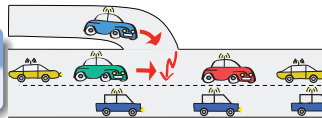
stochastic hybrid sys.

SHS = HS + stochastic



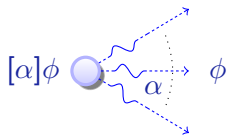
distributed hybrid sys.

DHS = HS + distributed



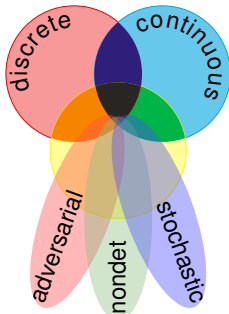
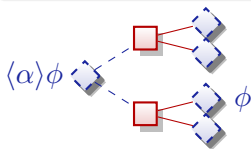
differential dynamic logic

$$d\mathcal{L} = DL + HP$$



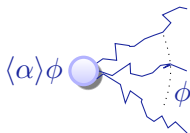
differential game logic

$$d\mathcal{G}\mathcal{L} = GL + HG$$



stochastic differential DL

$$Sd\mathcal{L} = DL + SHP$$

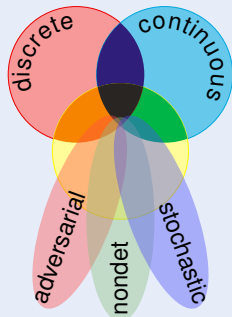


quantified differential DL

$$Qd\mathcal{L} = FOL + DL + QHP$$

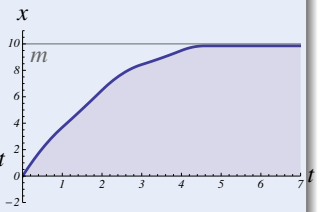
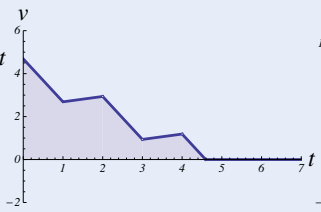
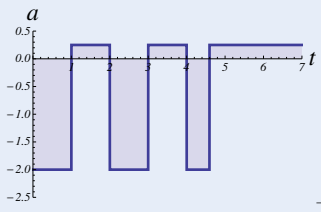
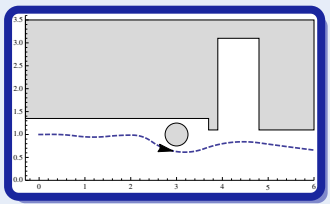
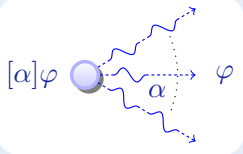
Dynamic Logics

- DL has been introduced for programs
Pratt'76,Harel,Kozen
- Its real calling are dynamical systems
- DL excels at providing simple+elegant
logical foundations for dynamical systems
- CPSs are multi-dynamical systems
- DL for CPS are multi-dynamical



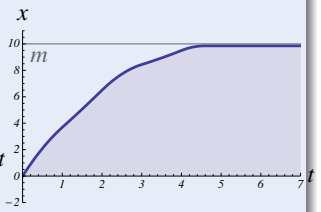
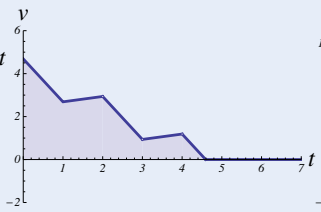
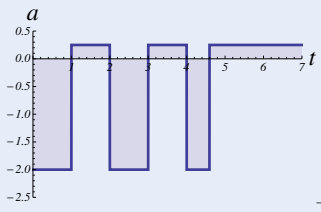
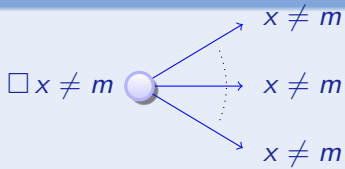
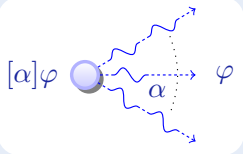
Concept (Differential Dynamic Logic)

(JAR'08, LICS'12)



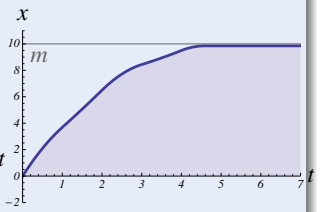
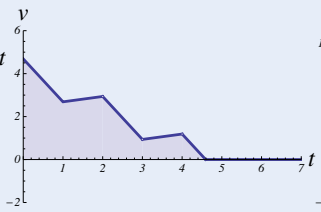
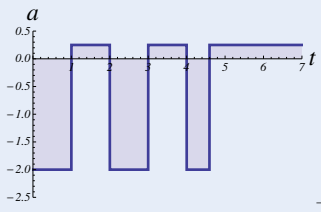
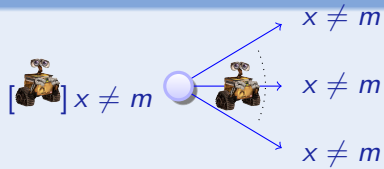
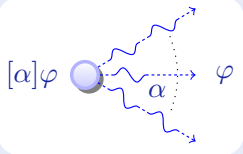
Concept (Differential Dynamic Logic)

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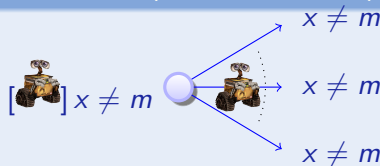
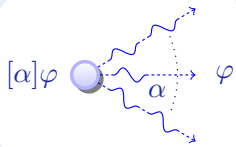
Concept (Differential Dynamic Logic)

(JAR'08, LICS'12)



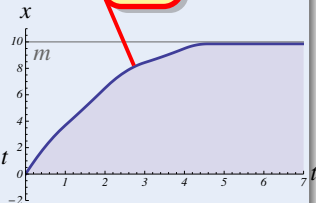
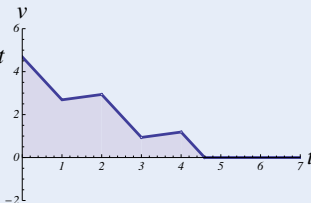
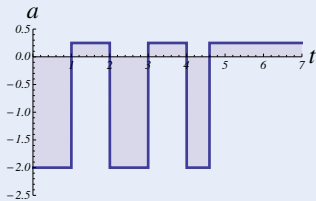
Concept (Differential Dynamic Logic)

(JAR'08, LICS'12)



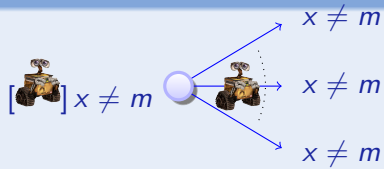
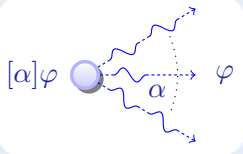
$$x' = v, v' = a$$

ODE



Concept (Differential Dynamic Logic)

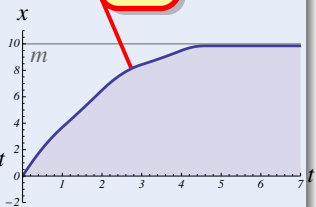
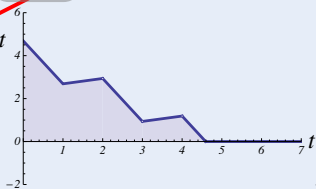
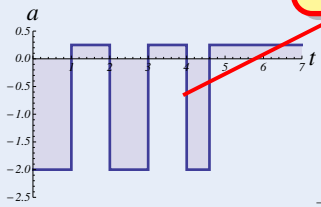
(JAR'08, LICS'12)



$$a := -b \quad x' = v, v' = a$$

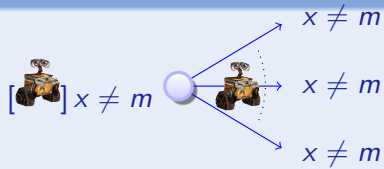
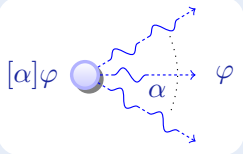
assign

ODE



Concept (Differential Dynamic Logic)

(JAR'08, LICS'12)

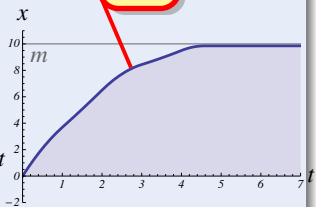
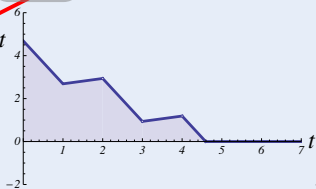
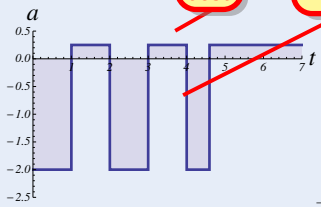


$$(\text{if}(\text{SB}(x, m)) a := -b) \quad x' = v, v' = a$$

test

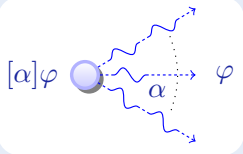
assign

ODE



Concept (Differential Dynamic Logic)

(JAR'08, LICS'12)



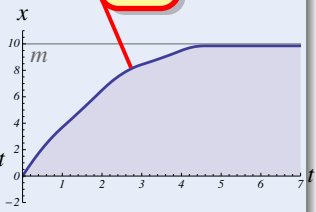
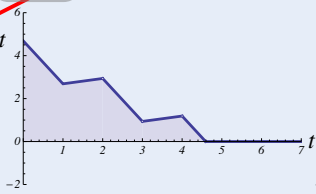
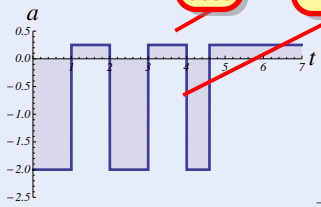
seq.
compose

$$(\text{if}(\text{SB}(x, m)) a := -b) ; x' = v, v' = a$$

test

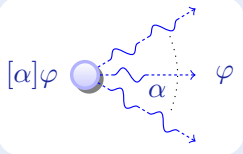
assign

ODE



Concept (Differential Dynamic Logic)

(JAR'08, LICS'12)



seq. compose

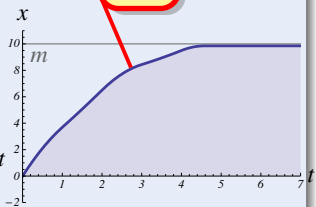
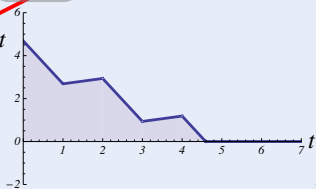
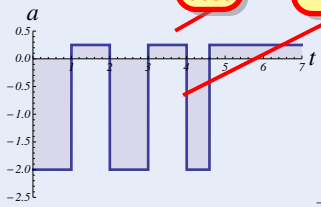
nondet. repeat

$$((\text{if}(\text{SB}(x, m)) a := -b) ; x' = v, v' = a)^*$$

test

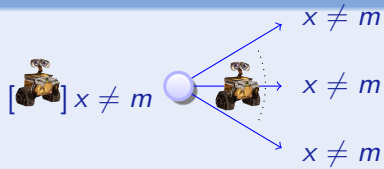
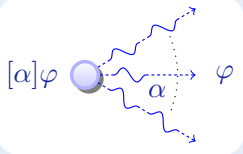
assign

ODE



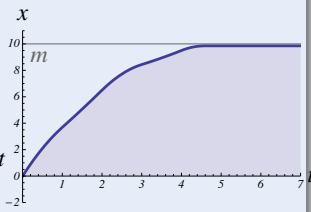
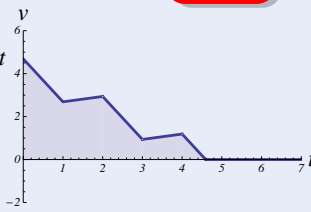
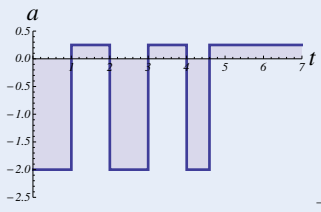
Concept (Differential Dynamic Logic)

(JAR'08, LICS'12)



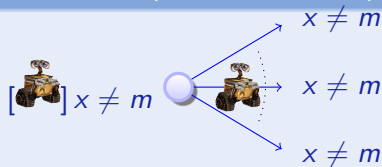
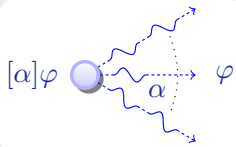
$$[((\text{if}(\text{SB}(x, m)) \ a := -b) ; \ x' = v, v' = a)^*] \underbrace{x \neq m}_{\text{post}}$$

all runs



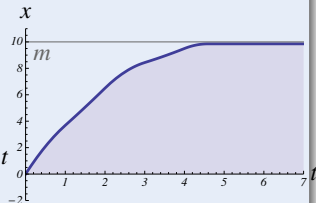
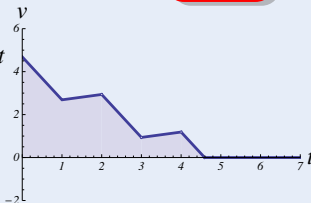
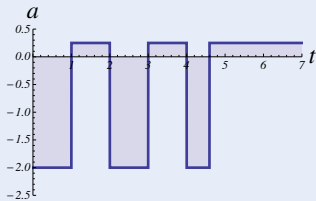
Concept (Differential Dynamic Logic)

(JAR'08, LICS'12)



$$\underbrace{x \neq m \wedge b > 0}_{\text{init}} \rightarrow \left[\left(\text{if}(\text{SB}(x, m)) \ a := -b \ ; \ x' = v, v' = a \right)^* \right] \underbrace{x \neq m}_{\text{post}}$$

all runs





Definition (Hybrid program α)

$$x := f(x) \mid ?Q \mid x' = f(x) \ \& \ Q \mid \alpha \cup \beta \mid \alpha; \beta \mid \alpha^*$$

Definition (d \mathcal{L} Formula P)

$$e \geq \tilde{e} \mid \neg P \mid P \wedge Q \mid \forall x P \mid \exists x P \mid [\alpha]P \mid \langle \alpha \rangle P$$



Differential Dynamic Logic dL: Syntax

Discrete Assign

Test Condition

Differential Equation

Nondet. Choice

Seq. Compose

Nondet. Repeat

Definition (Hybrid program α)

$x := f(x) \mid ?Q \mid x' = f(x) \ \& \ Q \mid \alpha \cup \beta \mid \alpha; \beta \mid \alpha^*$

Definition (dL Formula P)

$e \geq \tilde{e} \mid \neg P \mid P \wedge Q \mid \forall x P \mid \exists x P \mid [\alpha]P \mid \langle \alpha \rangle P$

All Reals

Some Reals

All Runs

Some Runs

JAR'08, LICS'12, JAR'17



$$[:=] \quad [x := e]P(x) \leftrightarrow P(e)$$

$$[?] \quad [?Q]P \leftrightarrow (Q \rightarrow P)$$

$$['] \quad [x' = f(x)]P \leftrightarrow \forall t \geq 0 [x := y(t)]P \quad (y'(t) = f(y))$$

$$[U] \quad [\alpha \cup \beta]P \leftrightarrow [\alpha]P \wedge [\beta]P$$

$$[;] \quad [\alpha; \beta]P \leftrightarrow [\alpha][\beta]P$$

$$[*] \quad [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P$$

$$K \quad [\alpha](P \rightarrow Q) \rightarrow ([\alpha]P \rightarrow [\alpha]Q)$$

$$I \quad [\alpha^*]P \leftrightarrow P \wedge [\alpha^*](P \rightarrow [\alpha]P)$$

$$C \quad [\alpha^*]\forall v > 0 (P(v) \rightarrow \langle \alpha \rangle P(v-1)) \rightarrow \forall v (P(v) \rightarrow \langle \alpha^* \rangle \exists v \leq 0 P(v))$$



Complete Proof Theory of Hybrid Systems

Theorem (Sound & Complete)

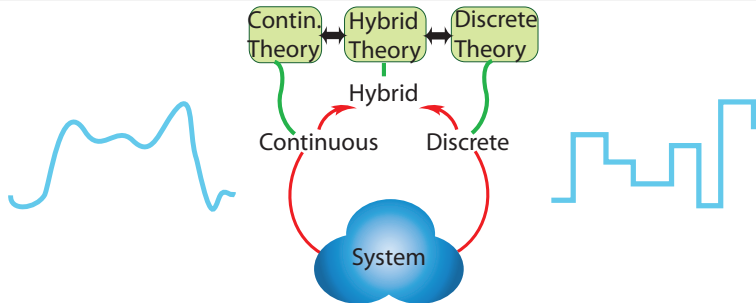
(JAR'08, LICS'12, JAR'17)

*dL calculus is a sound & complete axiomatization of hybrid systems relative to either differential equations **or** to discrete dynamics.*

▶ Proof 25pp

Corollary (Complete Proof-theoretical Bridge)

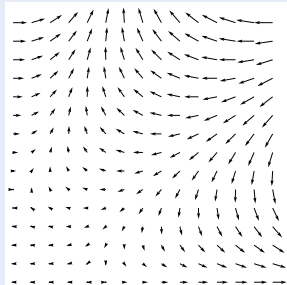
proving continuous = proving hybrid = proving discrete



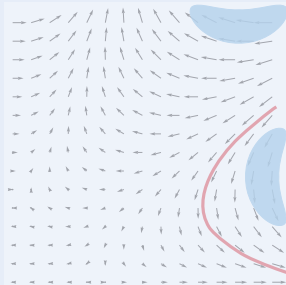


Differential Invariants for Differential Equations

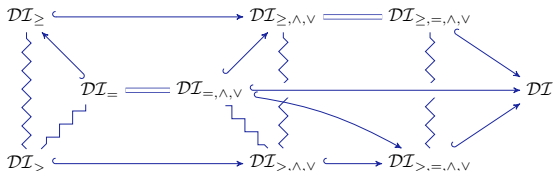
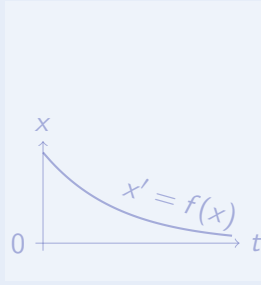
Differential Invariant



Differential Cut



Differential Ghost



Logic

Provability
theory

Math

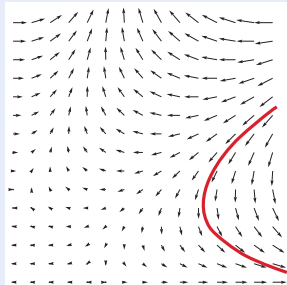
Characteristic
PDE

JLogComput'10, CAV'08, FMDS'09, LMCS'12, LICS'12, ITP'12, JAR'17

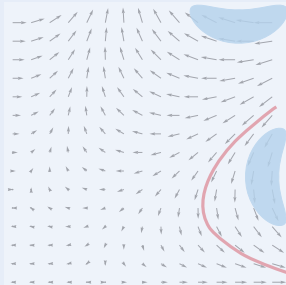


Differential Invariants for Differential Equations

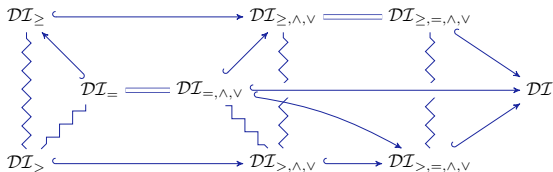
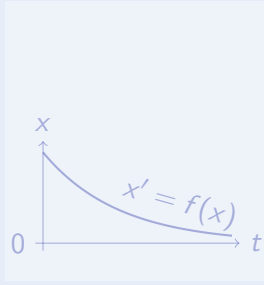
Differential Invariant



Differential Cut



Differential Ghost



Logic

Provability
theory

Math

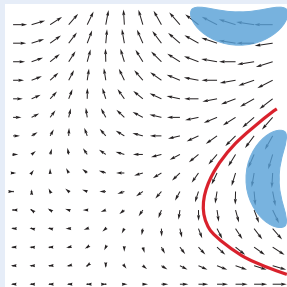
Characteristic
PDE

JLogComput'10, CAV'08, FMDS'09, LMCS'12, LICS'12, ITP'12, JAR'17

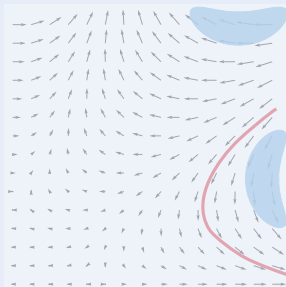


Differential Invariants for Differential Equations

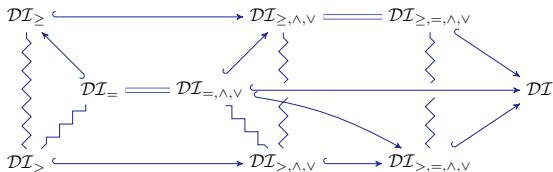
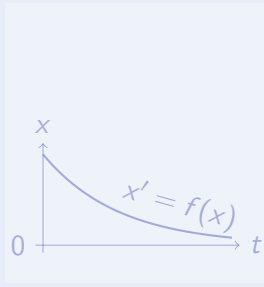
Differential Invariant



Differential Cut



Differential Ghost



Logic

Provability
theory

Math

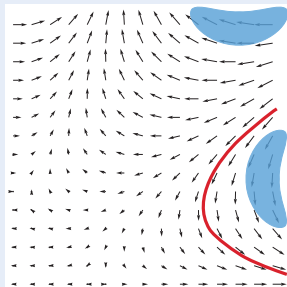
Characteristic
PDE

JLogComput'10, CAV'08, FMDS'09, LMCS'12, LICS'12, ITP'12, JAR'17

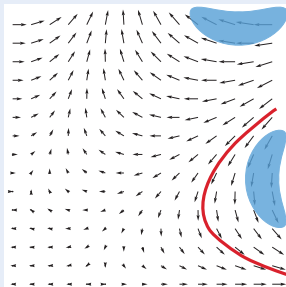


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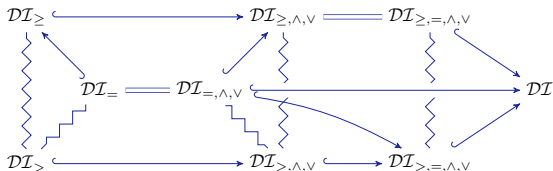
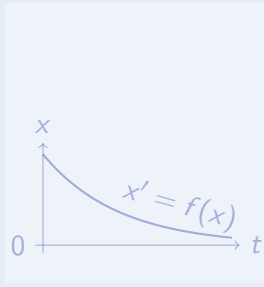
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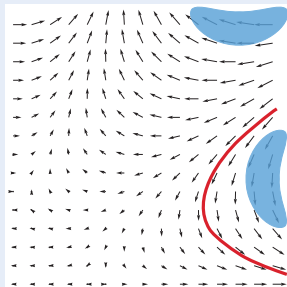
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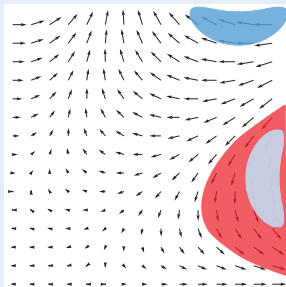


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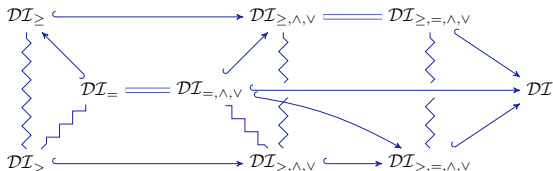
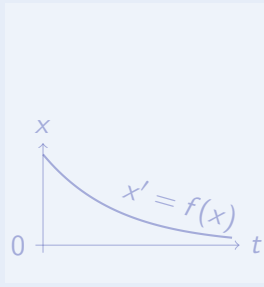
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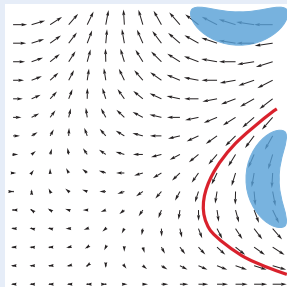
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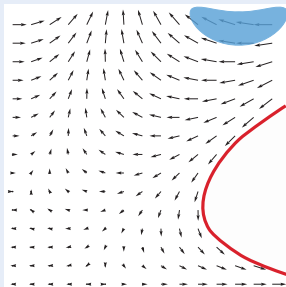


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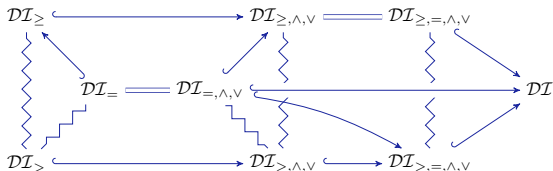
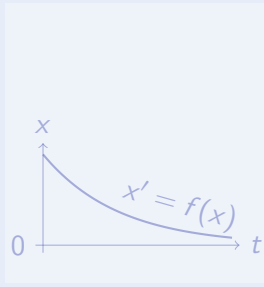
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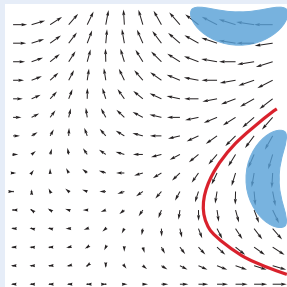
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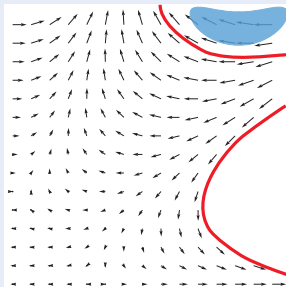


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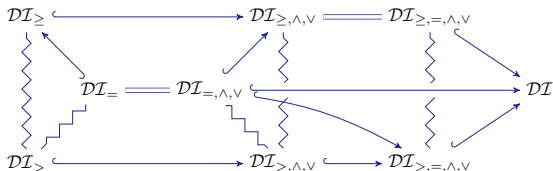
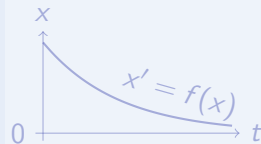
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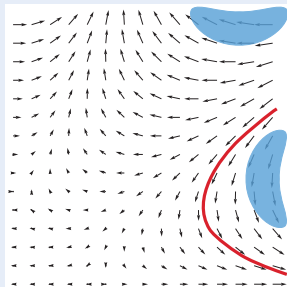
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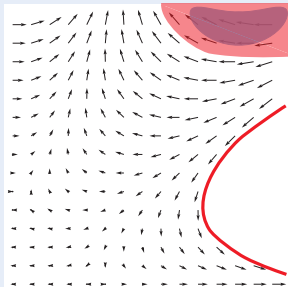
JLogComput'10, CAV'08, FMSD'09, LMCS'12, LICS'12, ITP'12, JAR'17

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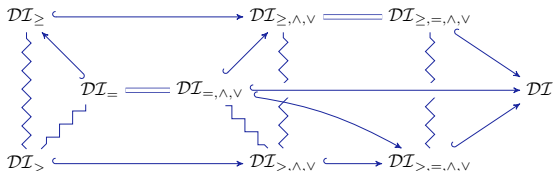
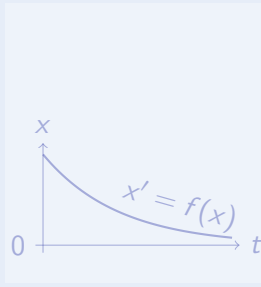
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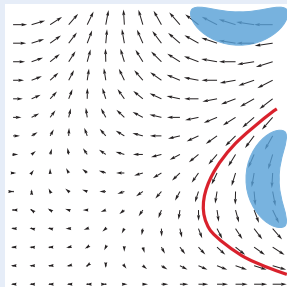
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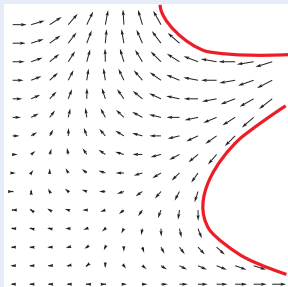


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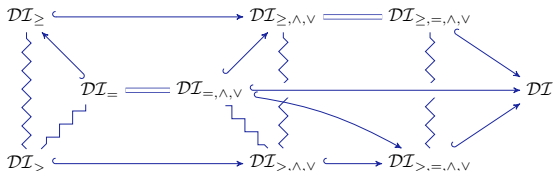
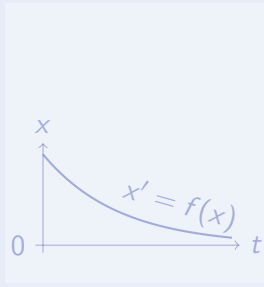
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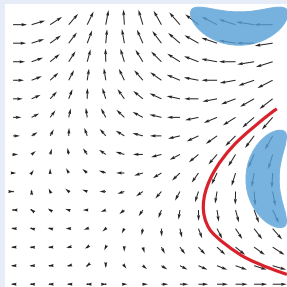
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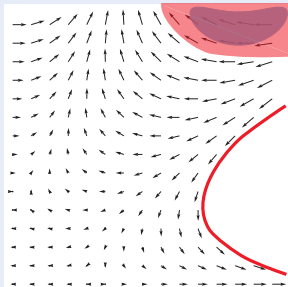
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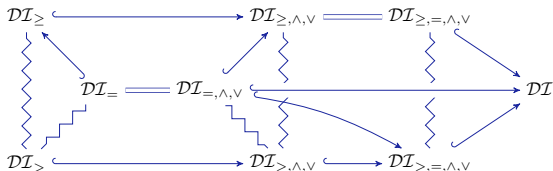
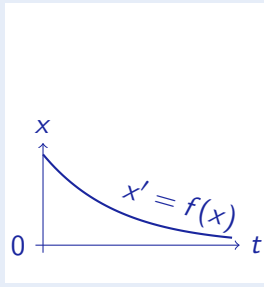
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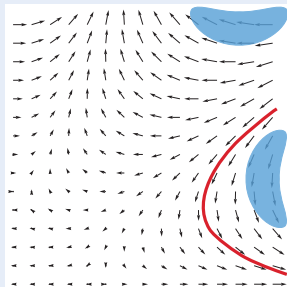
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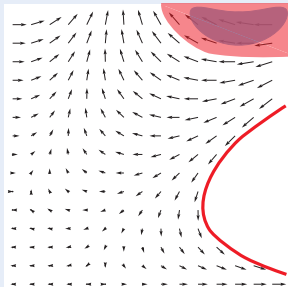
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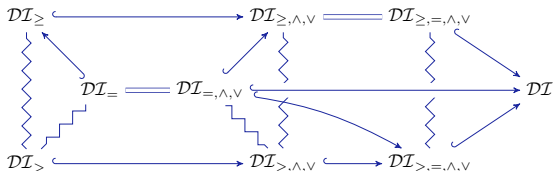
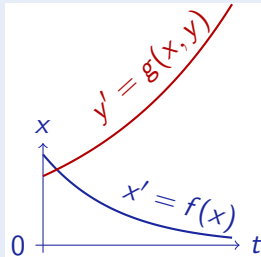
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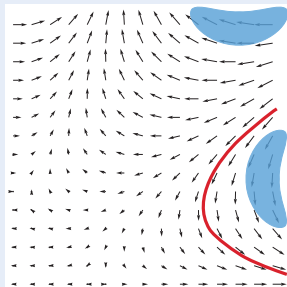
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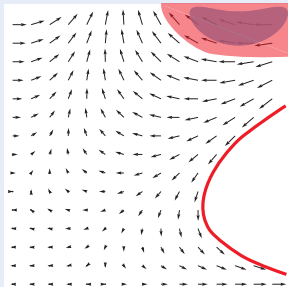


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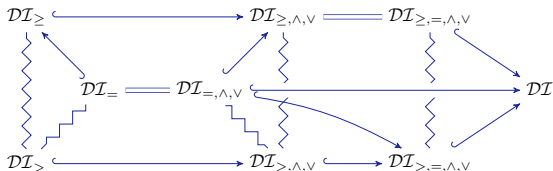
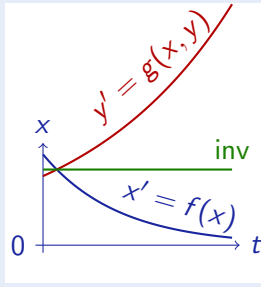
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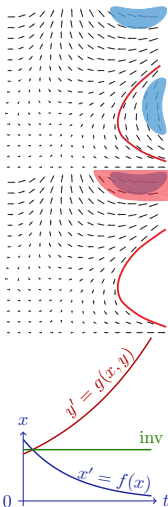
$$\frac{Q \vdash [x' := f(x)](P)'}{P \vdash [x' = f(x) \& Q]P}$$

Differential Cut

$$\frac{P \vdash [x' = f(x) \& Q]C \quad P \vdash [x' = f(x) \& Q \wedge C]P}{P \vdash [x' = f(x) \& Q]P}$$

Differential Ghost

$$\frac{P \leftrightarrow \exists y G \quad G \vdash [x' = f(x), y' = g(x, y) \& Q]G}{P \vdash [x' = f(x) \& Q]P}$$





Differential Invariant

$$\frac{Q \vdash [x' := f(x)](P)'}{P \vdash [x' = f(x) \& Q]P}$$

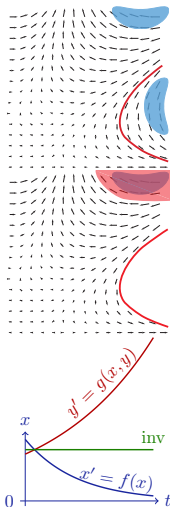
Differential Cut

$$\frac{P \vdash [x' = f(x) \& Q]C \quad P \vdash [x' = f(x) \& Q \wedge C]P}{P \vdash [x' = f(x) \& Q]P}$$

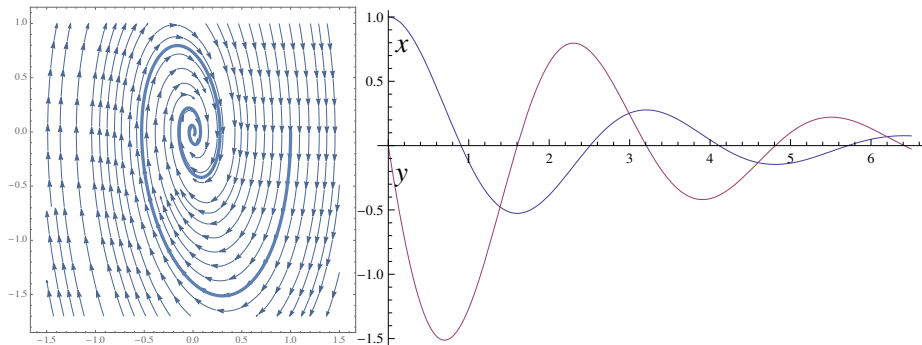
Differential Ghost

$$\frac{P \leftrightarrow \exists y G \quad G \vdash [x' = f(x), y' = g(x, y) \& Q]G}{P \vdash [x' = f(x) \& Q]P}$$

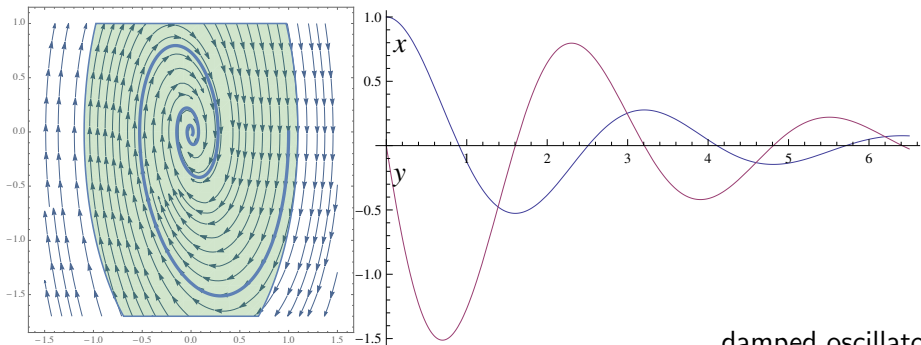
if new $y' = g(x, y)$ has a global solution



$$\omega^2 x^2 + y^2 \leq c^2 \vdash [x' = y, y' = -\omega^2 x - 2d\omega y \ \& \ \omega \geq 0 \wedge d \geq 0] \omega^2 x^2 + y^2 \leq c^2$$



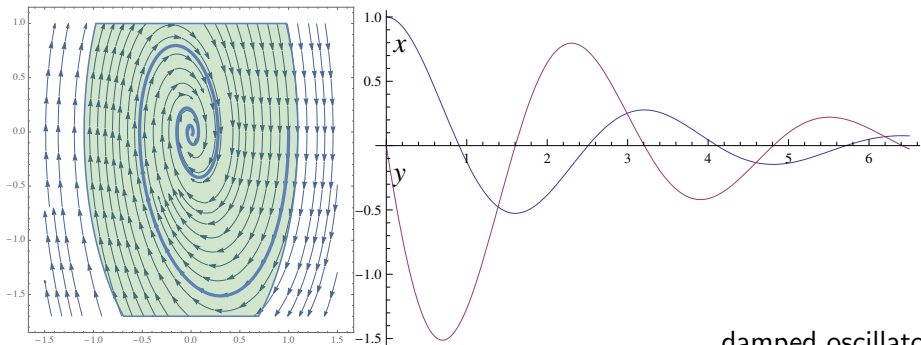
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damped oscillator

$$\omega \geq 0 \wedge d \geq 0 \vdash [x' := y][y' := -\omega^2 x - 2d\omega y] 2\omega^2 x x' + 2y y' \leq 0$$

$$\omega^2 x^2 + y^2 \leq c^2 \vdash [x' = y, y' = -\omega^2 x - 2d\omega y \ \& \ \omega \geq 0 \wedge d \geq 0] \omega^2 x^2 + y^2 \leq c^2$$

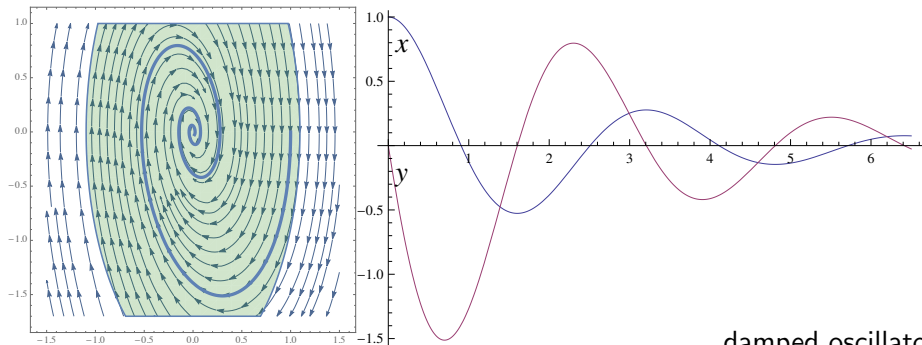


damped oscillator

$$\omega \geq 0 \wedge d \geq 0 \vdash 2\omega^2 xy + 2y(-\omega^2 x - 2d\omega y) \leq 0$$

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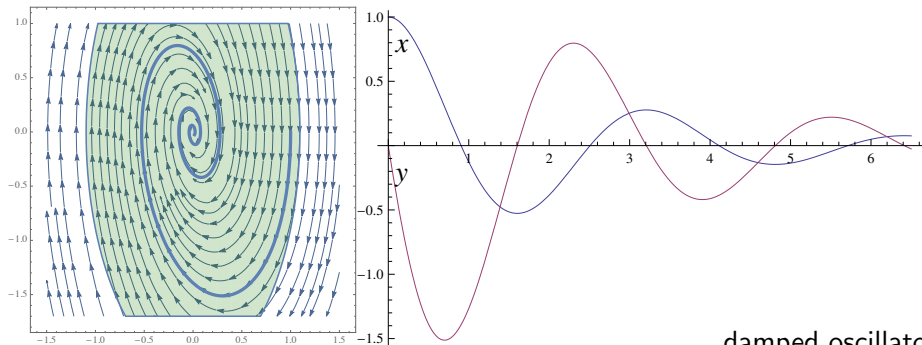
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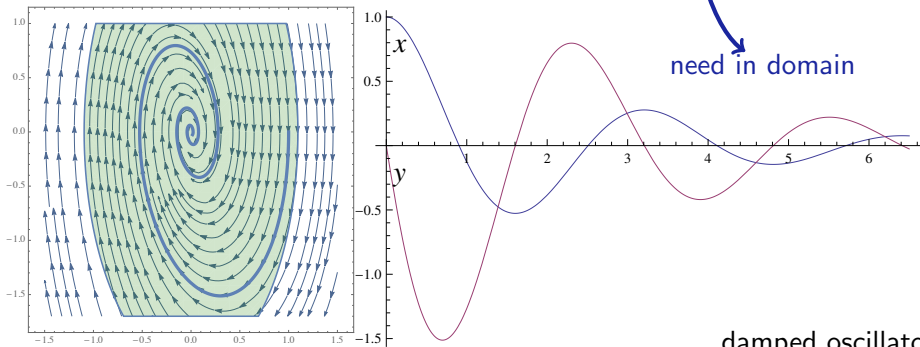
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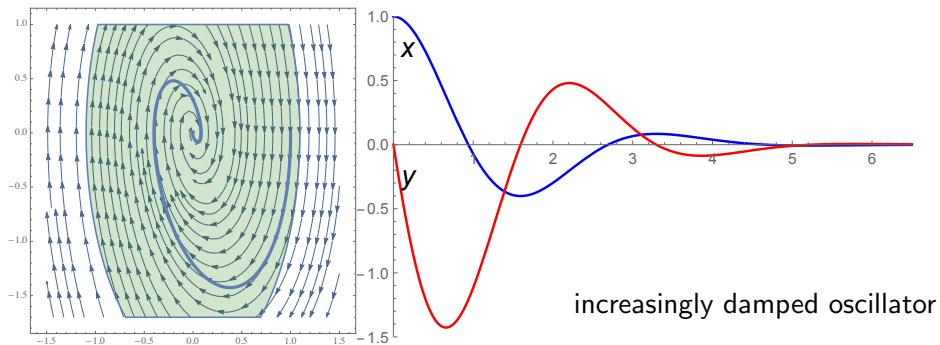
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damped oscillator

$$\omega^2 x^2 + y^2 \leq c^2 \vdash [x' = y, y' = -\omega^2 x - 2d\omega y, d' = 7 \ \& \ \omega \geq 0] \omega^2 x^2 + y^2 \leq c^2$$

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increasingly damped oscillator

$$\frac{\omega^2 x^2 + y^2 \leq c^2 \vdash [x' = y, y' = -\omega^2 x - 2d\omega y, d' = 7 \ \& \ \omega \geq 0 \ \& \ d \geq 0] \omega^2 x^2 + y^2 \leq c^2}{\omega^2 x^2 + y^2 \leq c^2 \vdash [x' = y, y' = -\omega^2 x - 2d\omega y, d' = 7 \ \& \ \omega \geq 0] \omega^2 x^2 + y^2 \leq c^2}$$

ask

$$\frac{}{d \geq 0 \vdash [x' = y, y' = -\omega^2 x - 2d\omega y, d' = 7 \ \& \ \omega \geq 0] d \geq 0}$$

increasingly damped oscillator

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$$\frac{}{\omega \geq 0 \vdash [d' := 7] d' \geq 0}$$

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DC

increasingly damped oscillator

$$\frac{\omega \geq 0 \wedge d \geq 0 \vdash [x' := y][y' := -\omega^2 x - 2d\omega y] 2\omega^2 x x' + 2y y' \leq 0}{\omega^2 x^2 + y^2 \leq c^2 \vdash [x' = y, y' = -\omega^2 x - 2d\omega y, d' = 7 \ \& \ \omega \geq 0 \wedge d \geq 0] \omega^2 x^2 + y^2 \leq c^2}$$

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increasingly damped oscillator

$$\omega \geq 0 \wedge d \geq 0 \vdash 2\omega^2 x y + 2y(-\omega^2 x - 2d\omega y) \leq 0$$

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increasingly damped oscillator

*

$$\omega \geq 0 \wedge d \geq 0 \vdash 2\omega^2xy + 2y(-\omega^2x - 2d\omega y) \leq 0$$

$$\omega \geq 0 \wedge d \geq 0 \vdash [x' := y][y' := -\omega^2x - 2d\omega y] 2\omega^2xx' + 2yy' \leq 0$$

$$\omega^2x^2 + y^2 \leq c^2 \vdash [x' = y, y' = -\omega^2x - 2d\omega y, d' = 7 \ \& \ \omega \geq 0 \wedge d \geq 0] \omega^2x^2 + y^2 \leq c^2$$

$$\omega^2x^2 + y^2 \leq c^2 \vdash [x' = y, y' = -\omega^2x - 2d\omega y, d' = 7 \ \& \ \omega \geq 0] \omega^2x^2 + y^2 \leq c^2$$

*

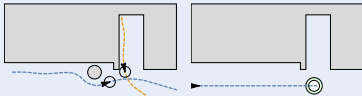
$$\omega \geq 0 \vdash 7 \geq 0$$

$$\omega \geq 0 \vdash [d' := 7] d' \geq 0$$

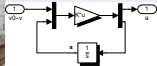
$$d \geq 0 \vdash [x' = y, y' = -\omega^2x - 2d\omega y, d' = 7 \ \& \ \omega \geq 0] d \geq 0$$

Could repeatedly diffcut in formulas to help the proof

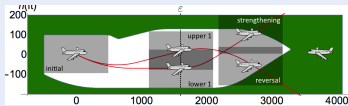
Obstacle Avoidance + Ground Navigation



Train Control Brakes



Airborne Collision Avoidance (ACAS X)



Ship Cooling



BOSCH SIEMENS



JOHNS HOPKINS
APPLIED PHYSICS LABORATORY



Ground Robot Obstacle Avoidance: Verify

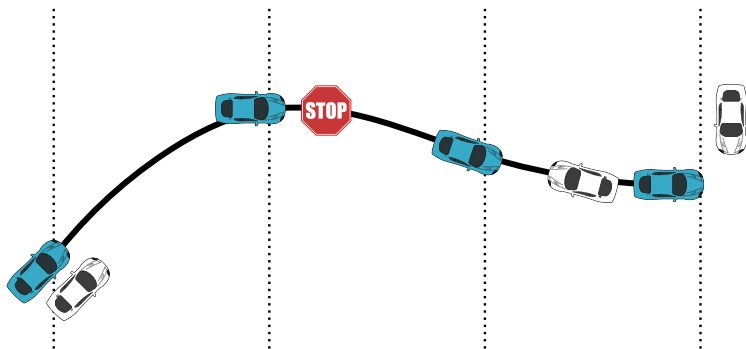
- Fundamental safety question for ground robot navigation
- When will which control decision avoid obstacles?
- Depends on safety objective, physical capabilities of robot + obstacle

Pass parking

Avoid/Follow

Head-on

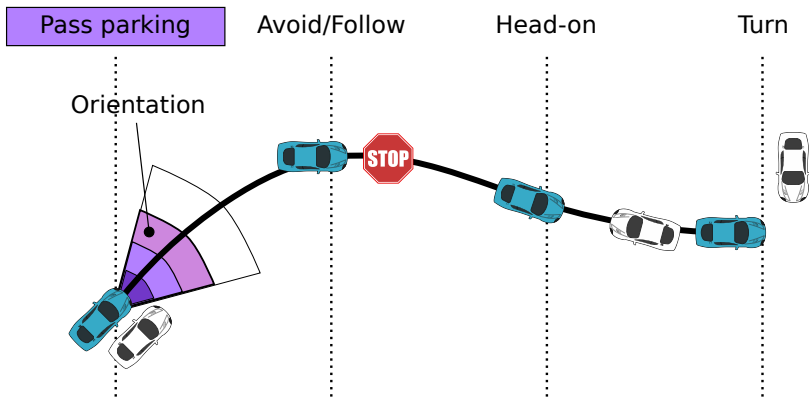
Turn



- 1 Identified safe region for each safety notion symbolically
- 2 Proved safety for hybrid systems ground robot model in KeYmaera X

Ground Robot Obstacle Avoidance: Verify

- Fundamental safety question for ground robot navigation
- When will which control decision avoid obstacles?
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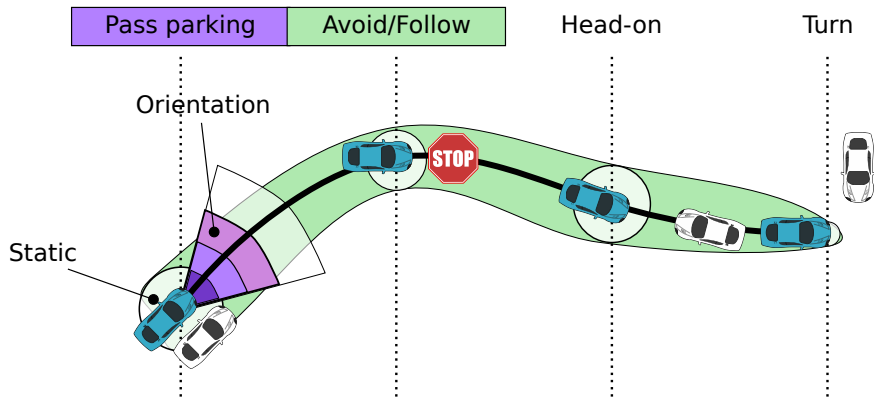


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Ground Robot Obstacle Avoidance: Verify

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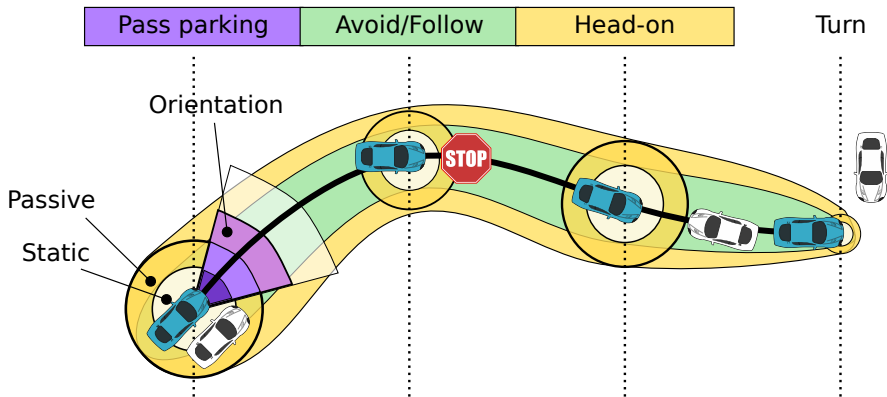


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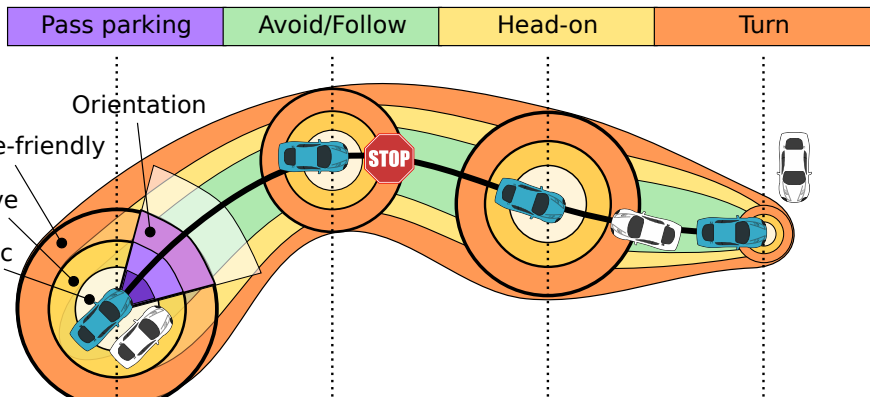


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Ground Robot Obstacle Avoidance: Verify

- Fundamental safety question for ground robot navigation
- When will which control decision avoid obstacles?
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- 2 Proved safety for hybrid systems ground robot model in KeYmaera X



Safety ▶

Invariant + Safe Control

static $\|p - o\|_\infty > \frac{s^2}{2b} + \left(\frac{A}{b} + 1\right) \left(\frac{A}{2}\varepsilon^2 + \varepsilon s\right)$

passive $s \neq 0 \rightarrow \|p - o\|_\infty > \frac{s^2}{2b} + V\frac{s}{b} + \left(\frac{A}{b} + 1\right) \left(\frac{A}{2}\varepsilon^2 + \varepsilon(s + V)\right)$

+ sensor $\|\hat{p} - o\|_\infty > \frac{s^2}{2b} + V\frac{s}{b} + \left(\frac{A}{b} + 1\right) \left(\frac{A}{2}\varepsilon^2 + \varepsilon(s + V)\right) + \Delta_p$

+ disturb $\|p - o\|_\infty > \frac{s^2}{2b\Delta_a} + V\frac{s}{b\Delta_a} + \left(\frac{A}{b\Delta_a} + 1\right) \left(\frac{A}{2}\varepsilon^2 + \varepsilon(s + V)\right)$

+ failure $\|\hat{p} - o\|_\infty > \frac{s^2}{2b} + V\frac{s}{b} + \left(\frac{A}{b} + 1\right) \left(\frac{A}{2}\varepsilon^2 + \varepsilon(v + V)\right) + \Delta_p + g\Delta$

friendly $\|p - o\|_\infty > \frac{s^2}{2b} + \frac{V^2}{2b_o} + V\left(\frac{s}{b} + \tau\right) + \left(\frac{A}{b} + 1\right) \left(\frac{A}{2}\varepsilon^2 + \varepsilon(s + V)\right)$



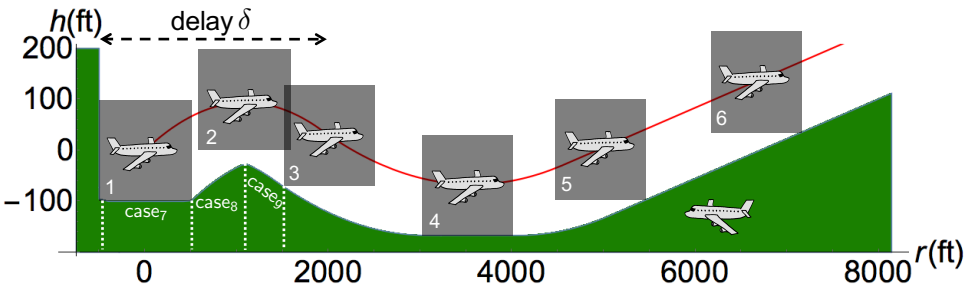
Safety	Invariant	+ Safe Control
static	$\ p - o\ _\infty > \frac{s^2}{2b}$	$+ \left(\frac{A}{b} + 1\right) \left(\frac{A}{2}\varepsilon^2 + \varepsilon s\right)$
passive	$s \neq 0 \rightarrow \ p - o\ _\infty > \frac{s^2}{2b}$	$+ V\frac{s}{b} + \left(\frac{A}{b} + 1\right) \left(\frac{A}{2}\varepsilon^2 + \varepsilon(s + V)\right)$
+ sensor		$+ \Delta_p$
+ disturb.	$\ p - o\ _\infty > \frac{s^2}{2b\Delta_a} + V\frac{s}{b\Delta_a}$	$+ \left(\frac{A}{b\Delta_a} + 1\right) \left(\frac{A}{2}\varepsilon^2 + \varepsilon(s + V)\right)$
+ failure	$\ \hat{p} - o\ _\infty > \frac{s^2}{2b} + V\frac{s}{b}$	$+ \left(\frac{A}{b} + 1\right) \left(\frac{A}{2}\varepsilon^2 + \varepsilon(v + V)\right) + \Delta_p + g\Delta$
friendly	$\ p - o\ _\infty > \frac{s^2}{2b} + \frac{V^2}{2b_o} + V\left(\frac{s}{b} + \tau\right)$	$+ \left(\frac{A}{b} + 1\right) \left(\frac{A}{2}\varepsilon^2 + \varepsilon(s + V)\right)$

Question
 How to find and justify constraints? Proof!



Airborne Collision Avoidance System ACAS X: Verify

- Developed by the FAA to replace current TCAS in aircraft
- Approximately optimizes Markov Decision Process on a grid
- Advisory from lookup tables with numerous 5D interpolation regions

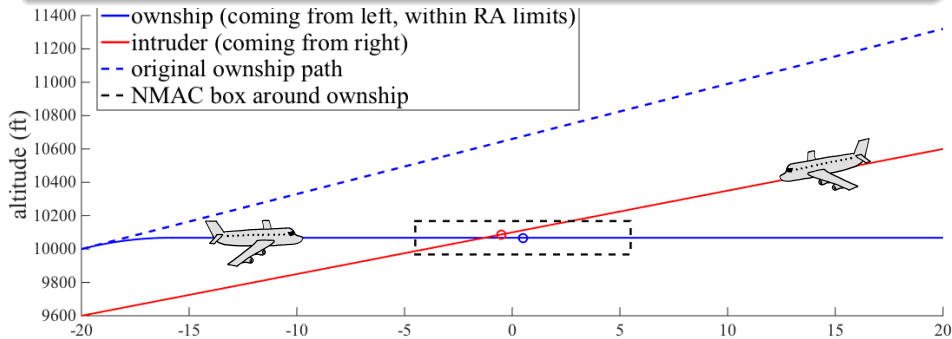


- 1 Identified safe region for each advisory symbolically
- 2 Proved safety for hybrid systems flight model in KeYmaera X

TACAS'15, EMSOFT'15, STTT'17



ACAS X table comparison shows safe advisory in 97.7% of the 648,591,384,375 states compared (15,160,434,734 counterexamples).

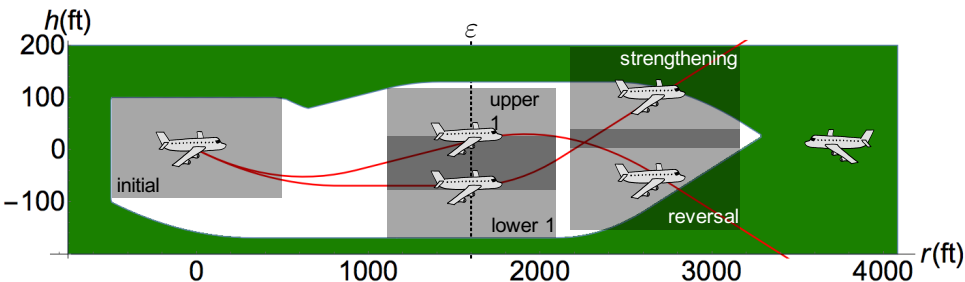


ACAS X issues DNC advisory, which induces collision unless corrected



Airborne Collision Avoidance System ACAS X: Refine

- Conservative, so too many counterexamples
- Settle for: safe for a little while, with safe future advisory possibility
- Safeable advisory: a subsequent advisory can safely avoid collision

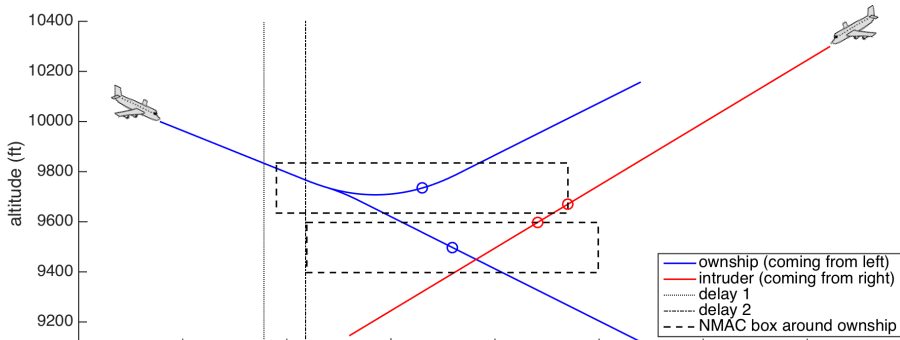


- 1 Identified safeable region for each advisory symbolically
- 2 Proved safety for hybrid systems flight model in KeYmaera X



ACAS X table comparison shows safeable advisory in more of the 648,591,384,375 states compared ($\approx 899 \cdot 10^6$ counterexamples).

**Counterexample: Action Issued = Maintain
Followed by Most Extreme Up/Down-sense Advisory Available**

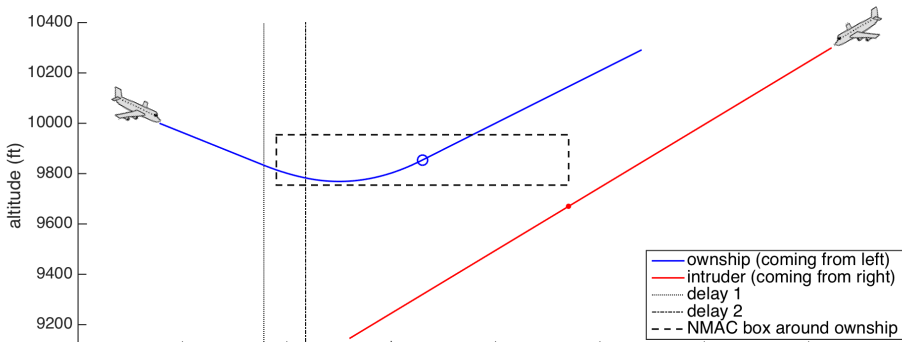


ACAS X issues Maintain advisory instead of CL1500



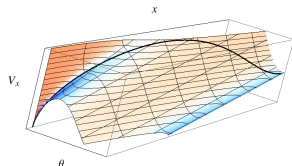
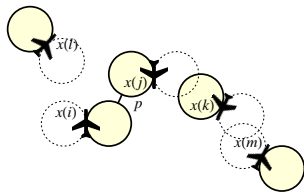
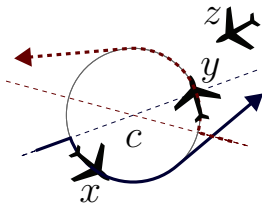
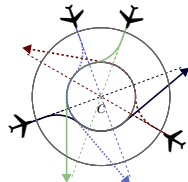
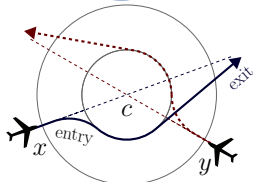
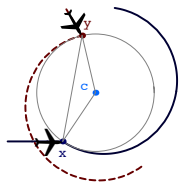
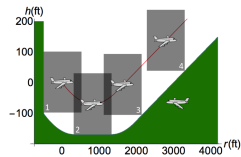
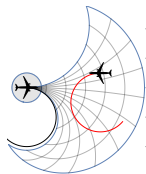
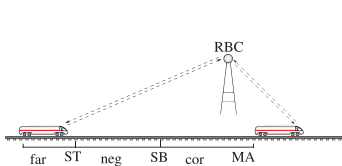
ACAS X table comparison shows safeable advisory in more of the 648,591,384,375 states compared ($\approx 899 \cdot 10^6$ counterexamples).

**Safe Version: Action Issued = CL1500
Followed by Most Extreme Up/Down-sense Available**



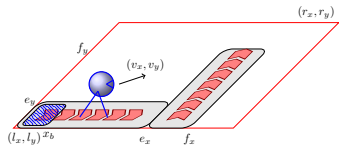
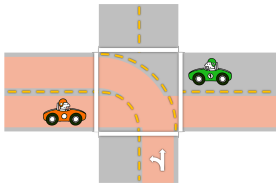
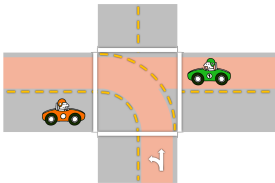
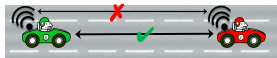
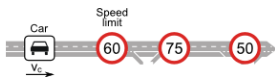
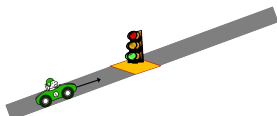
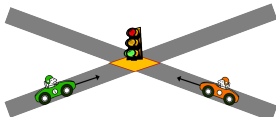
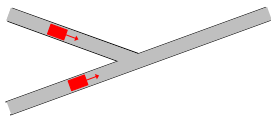
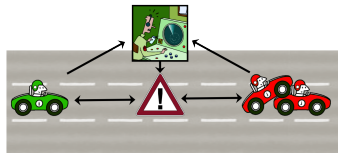
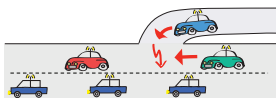
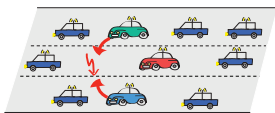
ACAS X issues Maintain advisory instead of CL1500

Verified CPS Applications: Trains & Airplanes



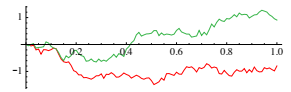
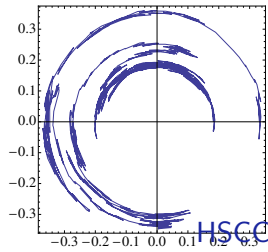
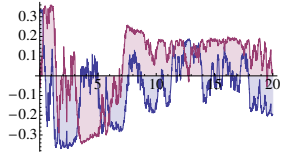
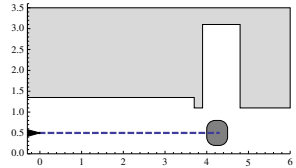
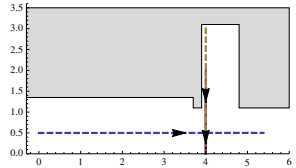
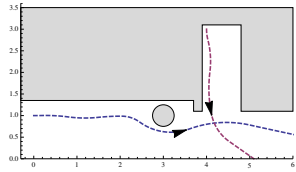
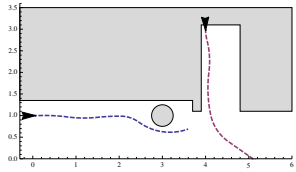
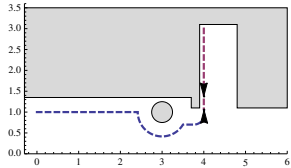
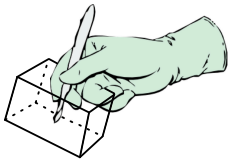
'09, JAIS'14, TACAS'15, EMSOFT'15, FM'09, HSCC'11, HSCC'13, TACAS'14, RSSRail'17

Verified CPS Applications: Cars

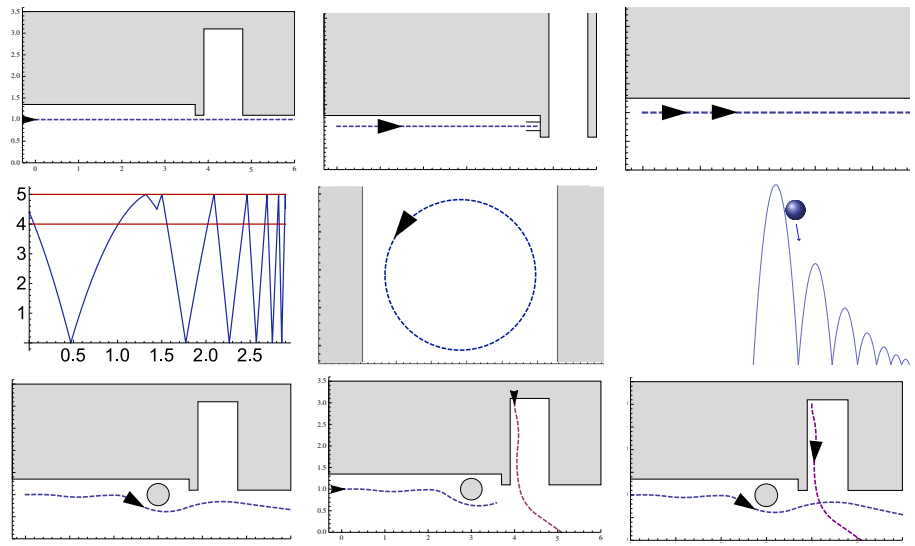


FM'11, LMCS'12, ICCPS'12, ITSC'11, ITSC'13, IJCAR'12

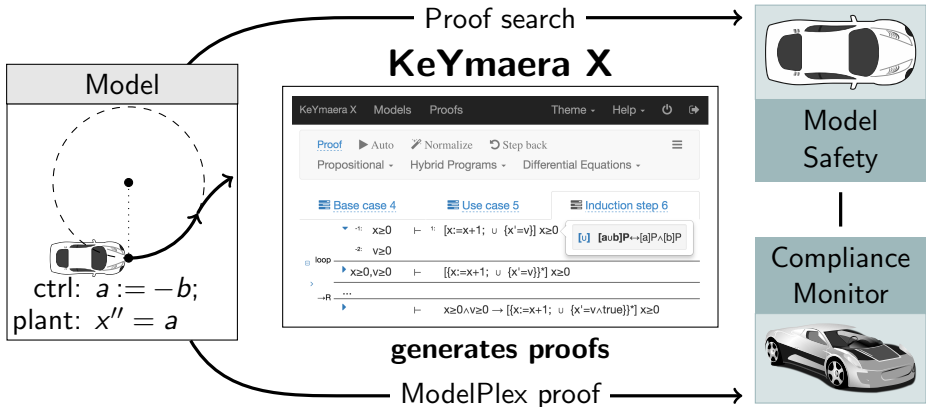
Verified CPS Applications: Robots



HSCC'13, RSS'13, CADE'12, IJRR'17



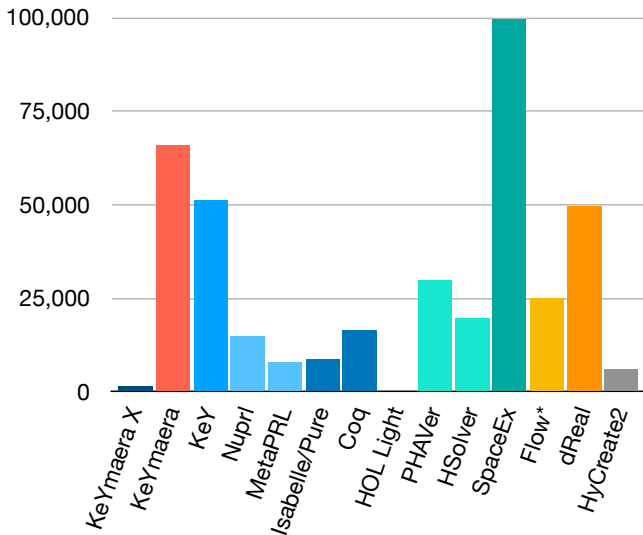
undergrads in *Foundations of Cyber-Physical Systems* course



Trustworthy
Uniform substitution
Sound & complete
Small core: 1700 LOC

Flexible
Proof automation
Interactive UI
Programmable

Customizable
Scala+Java API
Command line
REST API



Disclaimer: Self-reported estimates of the soundness-critical lines of code + rules



Acknowledgments

Students and postdocs of the Logical Systems Lab at Carnegie Mellon
Brandon Bohrer, Nathan Fulton, Sarah Loos, João Martins, Yong Kiam Tan
Khalil Ghorbal, Jean-Baptiste Jeannin, Stefan Mitsch



BOSCH **SIEMENS**

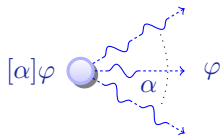
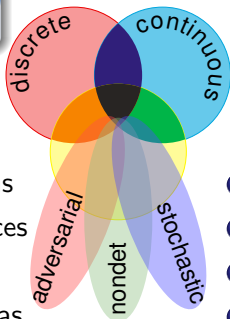


Logical foundations make a big difference for CPS, and vice versa

differential dynamic logic

$$d\mathcal{L} = DL + HP$$

- Strong analytic foundations
- Practical reasoning advances
- Significant applications
- Catalyze many science areas



- 1 Multi-dynamical systems
- 2 Combine simple dynamics
- 3 Tame complexity
- 4 www.keymaeraX.org

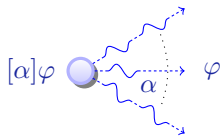
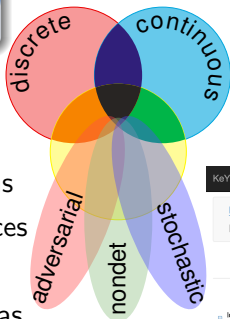
Numerous wonders remain to be discovered

Logical foundations make a big difference for CPS, and vice versa

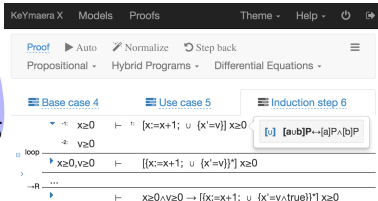
differential dynamic logic

$$d\mathcal{L} = DL + HP$$

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KeYmaera X

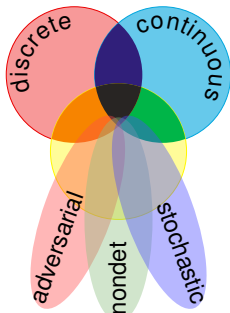


Numerous wonders remain to be discovered

Numerous wonders remain to be discovered

- Scalable continuous stochastics CADE'11
- Concurrent CPS
- Real arithmetic: Scalable and verified CADE'09
- Verified CPS implementations, ModelPlex FMDS'16
- Correct CPS execution
- CPS-conducive tactic languages+libraries ITP'17
- Tactics exploiting CPS structure/linearity/...
- Invariant generation FMDS'09 TACAS'14
- Tactics & proofs for reachable set computations
- Parallel proof search & disprovers
- Correct model transformation FM'14
- Inspiring applications

CPSs deserve proofs as safety evidence!



Overview

Cyber-physical systems (CPS) combine cyber capabilities, such as computation or communication, with physical capabilities, such as motion or other physical processes. Cars, aircraft, and robots are prime examples. Besides their novel physicality, it is often hard to determine by logical computational control algorithms. Designing these algorithms is challenging due to their tight coupling with physical behavior, which is often hard to model algorithmically. This book is the first book for safety-critical CPS. This textbook teaches undergraduate students for core principles behind CPSs. It shows them how to design models and analyze, identify safety specifications and control properties, understand abstraction and system architecture, design by synthesis, reason algorithmically about CPS models, verify CPS models of aggregate scale, and develop an intuition for operational effects. The book is supported with detailed lecture notes, lecture videos, homework assignments, and lab assignments.

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- 1.2. Dynamic Equations and Dynamic Systems
- 1.3. Control of Linear and Nonlinear Systems
- 1.4. Control of Rotational Systems and Robotics

Part II - Differential Equations Analysis

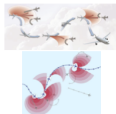
- 2.1. Differential Equations and Differential Systems
- 2.2. Differential Equations and Control
- 2.3. Differential Equations and Hybrid Systems

Part III - Mathematical Cyber-Physical Systems

- 3.1. Hybrid Systems and Control
- 3.2. Modeling Hybrid and Robotic Systems
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- 3.4. Modeling and Control of Hybrid Systems

Part IV - Computational Cyber-Physical Systems

- 4.1. Systems and Control Algorithms
- 4.2. Hybrid Models and Hybrid Control
- 4.3. Hybrid Verification and Model Checking
- 4.4. Hybrid Verification and Model Checking



Comments

"This excellent textbook teaches design and analysis of cyber-physical systems with a logical and computational way of thinking. The presentation is exemplary for finding the right balance between rigorous mathematical formalization and illustrative case studies aimed at providing intuition to system design."
 [Rajeev Alur, University of Pennsylvania]

"The author has developed major important tools for the design and control of these cyber-physical systems that increasingly shape our lives. This book is a 'must' for anyone interested, engineer, and mathematicians designing cyber-physical systems."
 [Andrzej Nowak, Cornell University]

"This book provides a wonderful introduction to cyber-physical systems, covering fundamental concepts from computer science and control theory from the perspective of formal logic. The material is brought to life through many detailed examples, illustrations, and exercises. A wealth of background material is provided in the text and in an appendix for each chapter, which makes the book well-organized and accessible to university students of all levels."
 [Srinivas Aravamudan, Université Grenoble Alpes]



Logical Analysis of Hybrid Systems

Proving Theorems for Complex Dynamics



Definition (Hybrid program semantics)

($\llbracket \cdot \rrbracket : \text{HP} \rightarrow \wp(\mathcal{S} \times \mathcal{S})$)

$$\llbracket x := e \rrbracket = \{(\omega, \nu) : \nu = \omega \text{ except } \llbracket x \rrbracket \nu = \llbracket e \rrbracket \omega\}$$

$$\llbracket ?Q \rrbracket = \{(\omega, \omega) : \omega \in \llbracket Q \rrbracket\}$$

$$\llbracket x' = f(x) \rrbracket = \{(\varphi(0), \varphi(r)) : \varphi \models x' = f(x) \text{ for some duration } r\}$$

$$\llbracket \alpha \cup \beta \rrbracket = \llbracket \alpha \rrbracket \cup \llbracket \beta \rrbracket$$

$$\llbracket \alpha; \beta \rrbracket = \llbracket \alpha \rrbracket \circ \llbracket \beta \rrbracket$$

$$\llbracket \alpha^* \rrbracket = \bigcup_{n \in \mathbb{N}} \llbracket \alpha^n \rrbracket$$

compositional semantics

Definition (d \mathcal{L} semantics)($\llbracket \cdot \rrbracket : \text{Fml} \rightarrow \wp(\mathcal{S})$)

$$\llbracket e \geq \tilde{e} \rrbracket = \{\omega : \llbracket e \rrbracket \omega \geq \llbracket \tilde{e} \rrbracket \omega\}$$

$$\llbracket \neg P \rrbracket = \llbracket P \rrbracket^c$$

$$\llbracket P \wedge Q \rrbracket = \llbracket P \rrbracket \cap \llbracket Q \rrbracket$$

$$\llbracket \langle \alpha \rangle P \rrbracket = \llbracket \alpha \rrbracket \circ \llbracket P \rrbracket = \{\omega : \nu \in \llbracket P \rrbracket \text{ for some } \nu : (\omega, \nu) \in \llbracket \alpha \rrbracket\}$$

$$\llbracket [\alpha] P \rrbracket = \llbracket \neg \langle \alpha \rangle \neg P \rrbracket = \{\omega : \nu \in \llbracket P \rrbracket \text{ for all } \nu : (\omega, \nu) \in \llbracket \alpha \rrbracket\}$$

$$\llbracket \exists x P \rrbracket = \{\omega : \omega \downarrow^r \in \llbracket P \rrbracket \text{ for some } r \in \mathbb{R}\}$$



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