

# Formal Verification of Train Control with Air Pressure Brakes

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Michael Golm<sup>2</sup>   André Platzer<sup>1</sup>

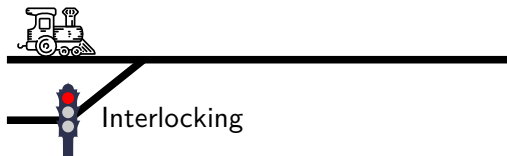
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<sup>2</sup>Siemens Corporate Technology, Princeton, NJ, USA

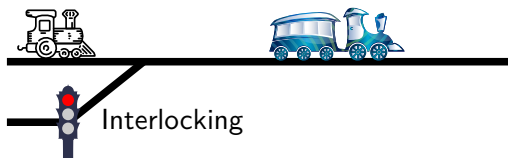
Reliability, Safety and Security of Railway Systems  
November 15, 2017



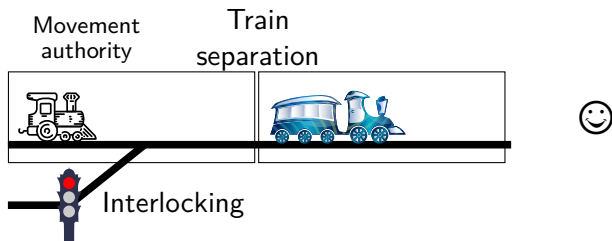
# Railroad Safety: Train Separation and Train Control



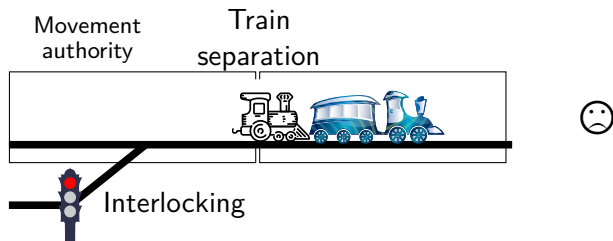
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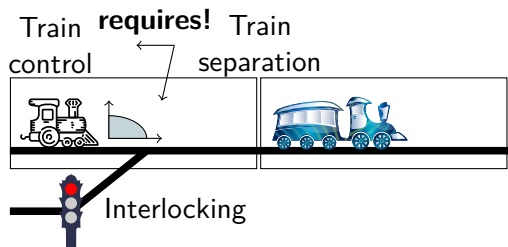
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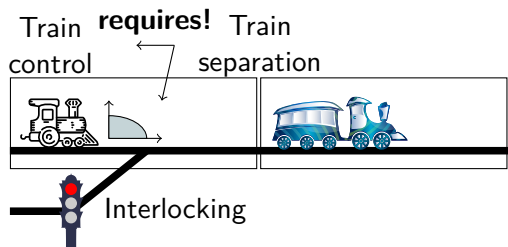
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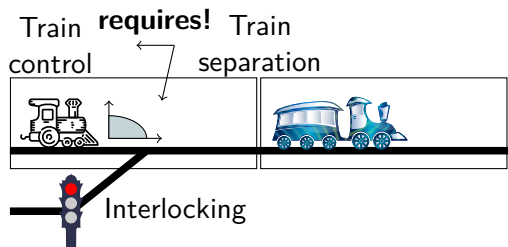
# Railroad Safety: Train Separation and Train Control



Design **provably safe train control** considering physical train motion

- Federal Railroad Administration (FRA): motion and brake models
- No overshoot
- Limited undershoot

# Railroad Safety: Train Separation and Train Control

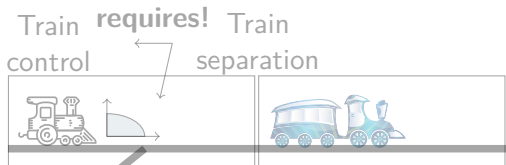


Design **provably safe train control** considering physical train motion

- Federal Railroad Administration (FRA): motion and brake models
- No overshoot
- Limited undershoot
- But underspecified control conditions**



# Railroad Safety: Train Separation and Train Control



## Approach

**Safe train separation  
requires verified train control and motion!**

Design **provably safe train control** considering physical train motion

Federal Railroad Administration (FRA): motion and brake models

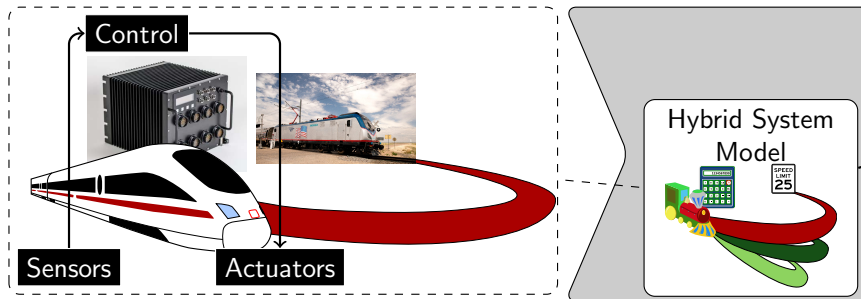
No overshoot

Limited undershoot

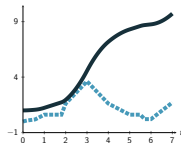
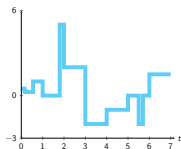
**But underspecified control conditions**

# Approach: Hybrid Systems Theorem Proving

Analyze the physical effect of software

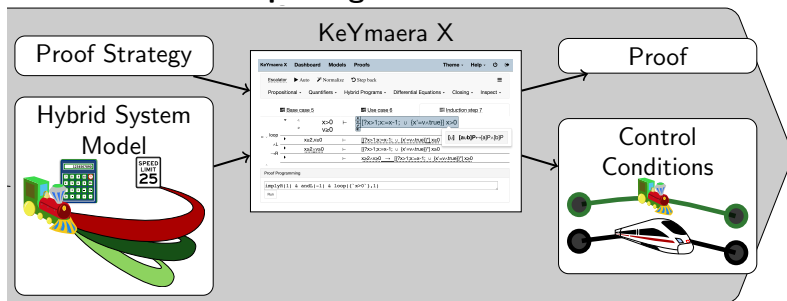


Discrete computation + continuous physics



# Approach: Hybrid Systems Theorem Proving

## Theorem proving ensures correct model



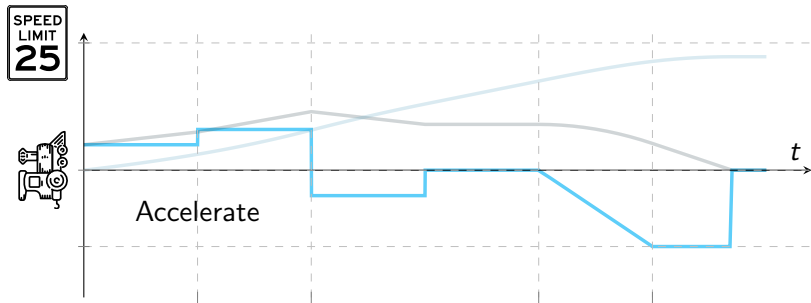
## Main results for

**Certification:** Proofs

System **architecture** and **implementation:** Models

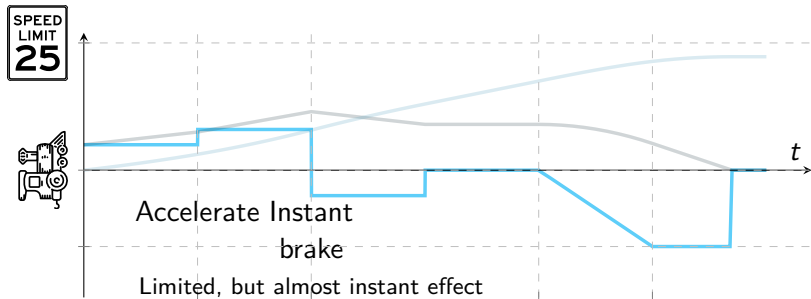
Control **engineering** and **testing:** Control conditions

# Train Motion and Brake Model

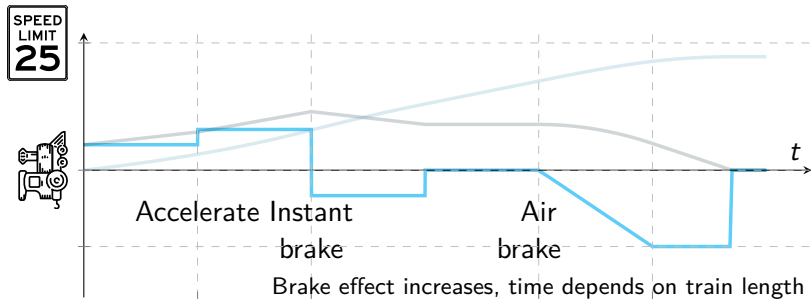


Brossaeu, J., Ede, B.M.: Development of an adaptive predictive braking enforcement algorithm. FRA/DOT/ORD-9/13 (2009)

# Train Motion and Brake Model

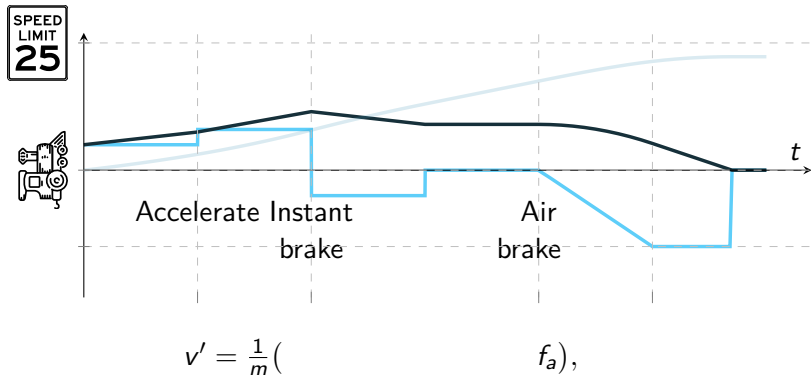


# Train Motion and Brake Model

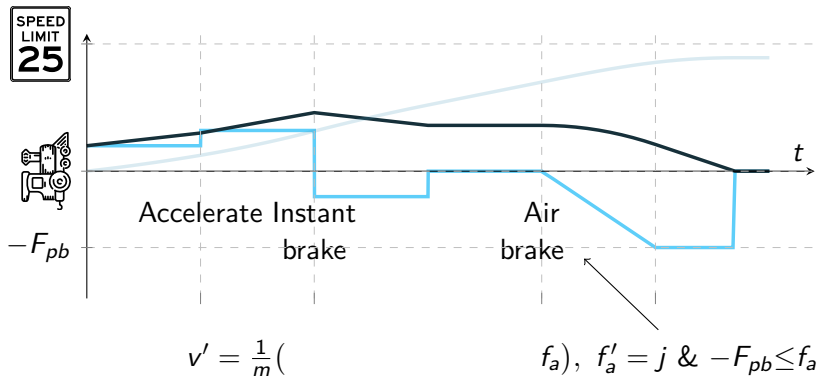


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# Train Motion and Brake Model

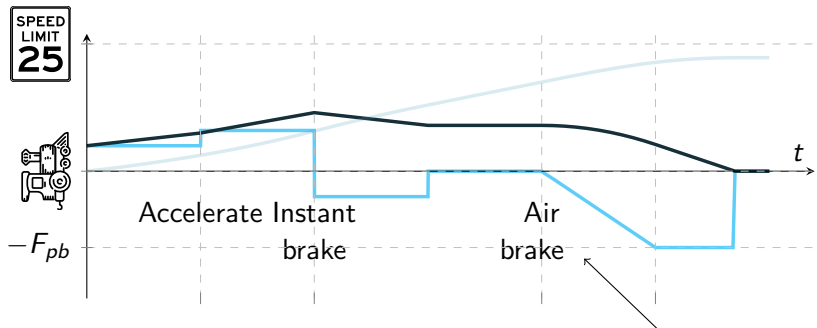


# Train Motion and Brake Model



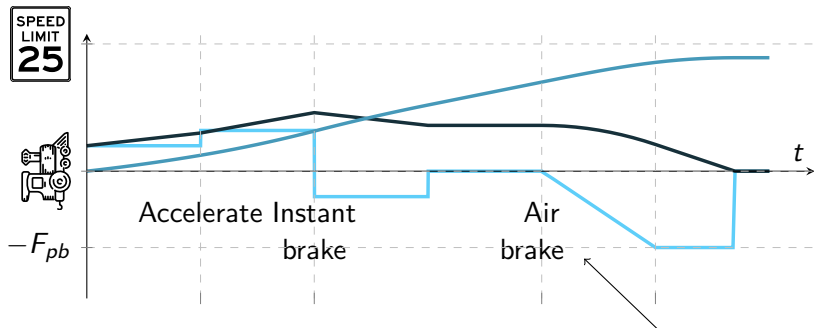


# Train Motion and Brake Model



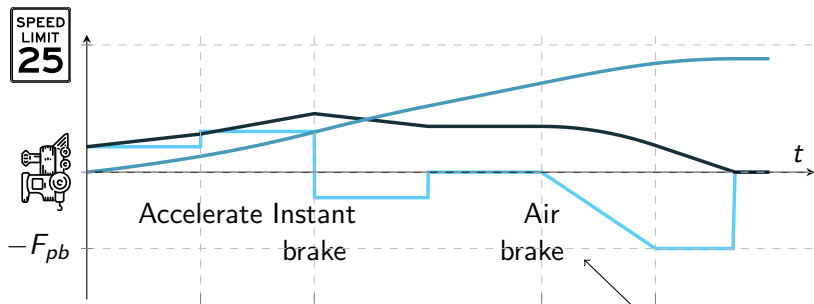
$$v' = \frac{1}{m} (-(F_g + F_r + F_c) + f_a), \quad f'_a = j \ \& \ -F_{pb} \leq f_a$$

# Train Motion and Brake Model



$$x' = v, \quad v' = \frac{1}{m}(- (F_g + F_r + F_c) + f_a), \quad f'_a = j \ \& \ -F_{pb} \leq f_a$$

# Train Motion and Brake Model



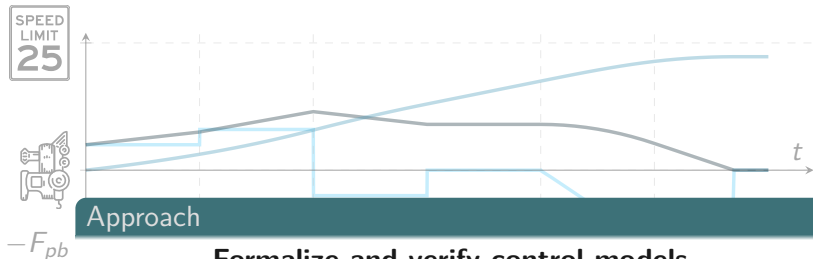
$$x' = v, \quad v' = \frac{1}{m} (-(F_g + F_r + F_c) + f_a), \quad f'_a = j \quad \& \quad -F_{pb} \leq f_a$$

Underspecified:

What are safe control choices?

How important is brake model fidelity?

# Train Motion and Brake Model



**Formalize and verify control models**  
**Analyze their brake engage points**

$$x' = v, \quad v' = \frac{1}{m}(- (F_g + F_r + F_c) + f_a), \quad f_a' = j \quad \& \quad -F_{pb} \leq f_a$$

Underspecified:

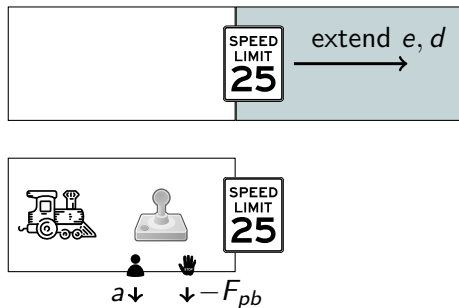
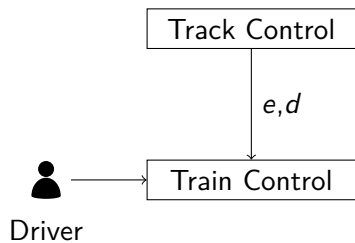
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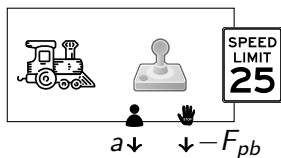
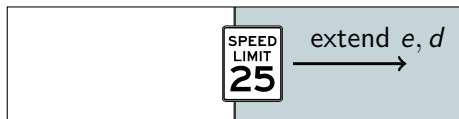
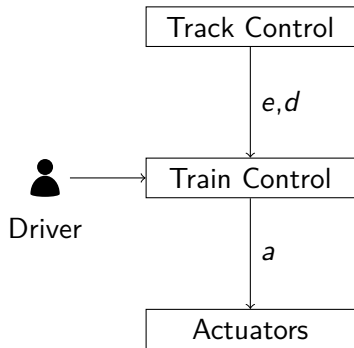
# Control Model





# Control Model

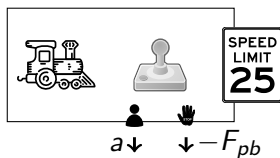
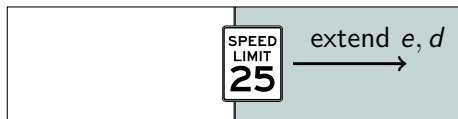
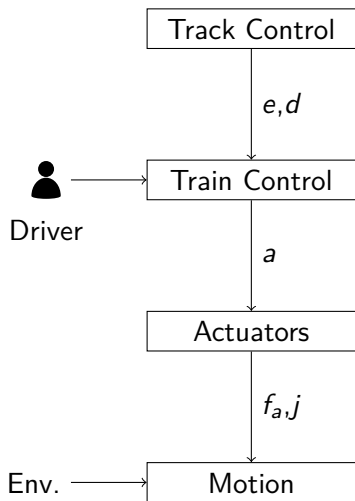




# Control Model



Delay  vs. air brake 

# Control Model



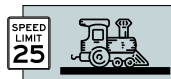
Delay  vs. air brake 



# Formal Verification with dL: No Overshoot

Correctness property: **respect the speed limit**

$$\text{safe} \equiv (z \geq e \rightarrow v \leq 25)$$



# Formal Verification with $d\mathcal{L}$ : No Overshoot

Correctness property: **respect the speed limit**

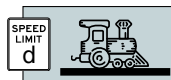
$$\text{safe} \equiv (z \geq e \rightarrow v \leq d)$$



# Formal Verification with $d\mathcal{L}$ : No Overshoot

Correctness property: **respect the speed limit**

$$\text{safe} \equiv (z \geq e \rightarrow v \leq d)$$

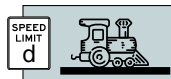


[]safe

# Formal Verification with $d\mathcal{L}$ : No Overshoot

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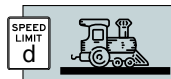


$$\left[ \left( \text{joystick}; \text{graph}; \text{train} \right)^* \right] \text{safe}$$

# Formal Verification with dL: No Overshoot

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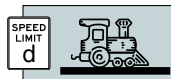


$$\left[ \left( \left[ \begin{array}{c} \text{SPEED} \\ \text{LIMIT} \\ d \end{array} \right] \cup \text{joystick}; \text{graph}; \text{train} \right)^* \right] \text{safe}$$

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Correctness property: **respect the speed limit**

$$\text{safe} \equiv (z \geq e \rightarrow v \leq d)$$



$$\text{init} \rightarrow \left[ \left( \begin{array}{c} \text{SPEED} \\ \text{LIMIT} \\ d \end{array} \cup \begin{array}{c} \text{joystick} \\ \text{; } \end{array} \begin{array}{c} \text{graph} \\ \text{; } \end{array} \begin{array}{c} \text{train} \end{array} \right)^* \right] \text{safe}$$

# Formal Verification with $d\mathcal{L}$ : No Overshoot

Correctness property: **respect the speed limit**

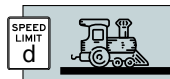
$$\text{safe} \equiv (z \geq e \rightarrow v \leq d)$$



$$\text{init} \rightarrow \left[ \left( \begin{array}{c} \text{SPEED} \\ \text{LIMIT} \\ d \end{array} \right) \cup \underbrace{\left( \text{person} \cup \text{STOP} \right)}_{\text{signal}} ; \begin{array}{c} \text{graph} \\ \text{train} \end{array} \right]^* \text{safe}$$

# Formal Verification with dL: No Overshoot

Correctness property: **respect the speed limit**



$$\text{safe} \equiv (z \geq e \rightarrow v \leq d)$$

$$\text{init} \rightarrow \left[ \left( \begin{array}{c} \text{SPEED LIMIT} \\ d \end{array} \cup \underbrace{\left( \text{person} \cup \text{STOP} \right)}_{\text{lock}} ; \begin{array}{c} \text{graph} \\ \text{train} \end{array} \right)^* \right] \text{safe}$$

$$\text{person} \equiv f_a := *; ? - F_{sb} \leq f_a \leq A; ? \left( e - z \geq \frac{(v^2 - d^2)m}{2F_{sb}} + \left( \frac{A}{F_{sb}} + 1 \right) \left( \frac{A}{2m} \varepsilon^2 + v\varepsilon \right) \right) \vee \text{slow}^- \vee \text{slow}^+ \vee \text{fast}^- \vee \text{fast}^+ \quad (1)$$

$$\text{slow}^- \equiv \neg \text{isFast}(v) \wedge f_a \leq 0 \wedge e - z \geq v\varepsilon + m\text{Slow}(v) \quad (3)$$

$$\text{slow}^+ \equiv [u := v + \frac{f_a \varepsilon}{m}] \left( \neg \text{isFast}(u) \wedge f_a \geq 0 \wedge e - z \geq v\varepsilon + \frac{f_a \varepsilon^2}{2m} + m\text{Slow}(u) \right) \quad (4)$$

$$\text{fast}^- \equiv \text{isFast}(v) \wedge f_a \leq 0 \wedge e - z \geq v\varepsilon + m\text{Fast}(v) \quad (5)$$

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$$\text{isFast}(v) \equiv v \geq \frac{F_{pb}^2}{2mJ} \quad m\text{Slow}(v) = \frac{2}{3} v \sqrt{2mv/J} \quad m\text{Fast}(v) = \frac{mv^2}{2F_{pb}} + \frac{vF_{pb}}{2J} - \frac{F_{pb}^3}{24mJ^2} \quad (7)$$



# Formal Verification with $d\mathcal{L}$ : No Overshoot

Correctness property: **respect the speed limit**



$$\text{safe} \equiv (z \geq e \rightarrow v \leq d)$$

$$\text{init} \rightarrow \left[ \left( \begin{array}{c} \text{SPEED LIMIT} \\ d \end{array} \cup \underbrace{(\text{person} \cup \text{STOP})}_{\text{driver}}; \text{graph}; \text{train} \right)^* \right] \text{safe}$$

## Result

**Driving with verified control conditions ensures safety**

How to find conditions: Proofs!

(1)

(2)

$$\text{slow}^- \equiv \neg \text{isFast}(v) \wedge f_a \leq 0 \wedge e - z \geq v\varepsilon + m\text{Slow}(v) \quad (3)$$

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# Systematically Derive Safe Control Conditions in $d\mathcal{L}$

**Partial model, unknown condition ?true**

$\left[ \text{person} ; ?\text{true} ; \text{train} \right] \text{safe}$

# Systematically Derive Safe Control Conditions in dL

**Partial model, unknown condition ?true**

$\left[ \text{person} ; ?\text{true} ; \text{train} \right] \text{safe}$

**Run and observe “parallel universe”**

$\left[ \text{person} ; \underbrace{?[\text{train}] \text{safe}} ; \text{train} \right] \text{safe}$

test for desired outcome



↓ copy&paste



# Systematically Derive Safe Control Conditions in dL

**Partial model, unknown condition ?true**

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test for desired outcome

Obviously true but not helpful for implementation

$\left[ \text{person} \right] \left( [\text{train}] \text{safe} \rightarrow [\text{train}] \text{safe} \right)$



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# Systematically Derive Safe Control Conditions in dL

**Partial model, unknown condition ?true**

$$\left[ \text{person} ; ?\text{true} ; \text{train} \right] \text{safe}$$

**Run and observe “parallel universe”**

$$\left[ \text{person} ; \underbrace{?[\text{train}] \text{safe}}_{\text{test for desired outcome}} ; \text{train} \right] \text{safe}$$

test for desired outcome

Obviously true but not helpful for implementation

$$\left[ \text{person} \right] \left( [\text{train}] \text{safe} \rightarrow [\text{train}] \text{safe} \right)$$

**Symbolically execute program with dL**

$$\left[ \text{person} \right] \left( \text{safe} \left( z + vt + \frac{f_a}{2m} t^2, v + \frac{f_a}{m} t \right) \rightarrow [\text{train}] \text{safe} \right)$$



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$$\text{train} \equiv \sum_{i=1}^n i \frac{n(n+1)}{2}$$

# Systematically Derive Safe Control Conditions in dL

**Partial model, unknown condition ?true**

$[ \text{person} ; ?\text{true} ; \text{train} ] \text{safe}$

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**Symbolically execute program with dL**

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**Use program effects as control conditions**

$[ \text{person} ; ?\text{safe}(z + vt + \frac{f_a}{2m} t^2, v + \frac{f_a}{m} t) ; \text{train} ] \text{safe}$



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$$\equiv \sum_{i=1}^n i \frac{n(n+1)}{2}$$

**Implementable**

# Systematically Derive Safe Control Conditions in dL

Partial model, unknown condition ?true

[  ; ?true ;  ] safe

Run and observe “parallel universe”

[  ; ?[  ] safe ;  ] safe





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Result

**Augment partial model  
with implementable control conditions  
derived by proof**

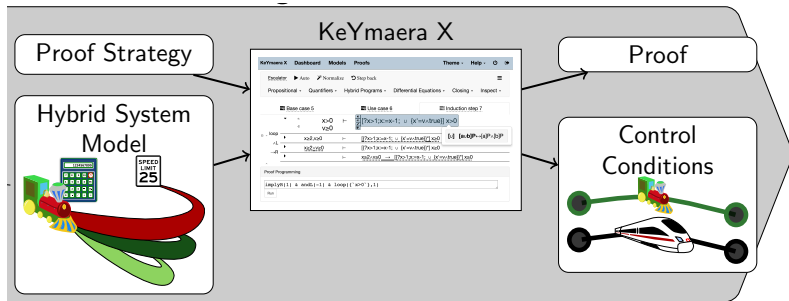
[  ] ( safe( $z + vt + \frac{f_a}{2m} t^2, v + \frac{f_a}{m} t$ ) → [  ] safe )

Use program effects as control conditions

[  ; ?safe( $z + vt + \frac{f_a}{2m} t^2, v + \frac{f_a}{m} t$ ) ;  ] safe

**Implementable**

$$\sum_{i=1}^n i = \frac{n(n+1)}{2}$$

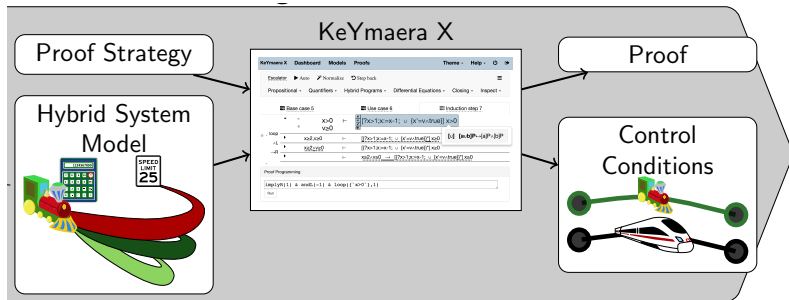


**No overshoot ✓**

Verified models: safely control brake delay and air brakes

Symbolic control conditions to select between free driving and braking





**No overshoot** ✓ + Limited undershoot

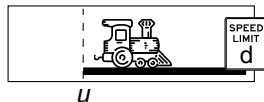
Verified models: safely control brake delay and air brakes

Symbolic control conditions to select between free driving and braking

# Formal Verification with $d\mathcal{L}$ : Limited Undershoot

Correctness: **when done braking, train is after undershoot limit**

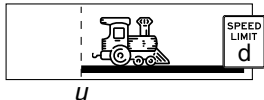
efficient  $\equiv (\text{brakesEngaged} \wedge v \leq d) \rightarrow (z \geq e - u)$



# Formal Verification with $d\mathcal{L}$ : Limited Undershoot

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Needs change in control priority

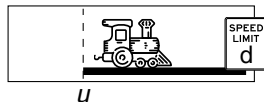
No overshoot: braking allowed but spoils efficiency

init  $\rightarrow \left[ \left( \left( \text{person} \cup \text{hand}_{\text{STOP}} \right) ; \text{signal} ; \text{train} \right)^* \right] \text{efficient}$

# Formal Verification with $d\mathcal{L}$ : Limited Undershoot

Correctness: **when done braking, train is after undershoot limit**

efficient  $\equiv (\text{brakesEngaged} \wedge v \leq d) \rightarrow (z \geq e - u)$

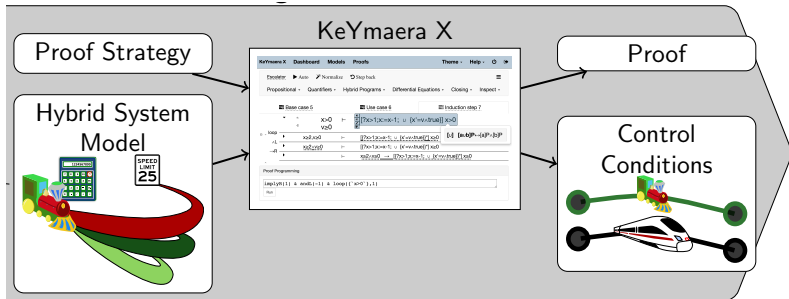


Needs change in control priority

No overshoot: braking allowed but spoils efficiency

Limited undershoot: only brake if necessary for safety

init  $\rightarrow \left[ \left( (\text{if } (\text{mustBrake}) \text{ STOP} \text{ else } \text{person}) ; \text{L} ; \text{train} \right)^* \right] \text{efficient}$



**No overshoot ✓ + Limited undershoot ✓**

Verified models: safely control brake delay and air brakes

Symbolic control conditions to select between free driving and braking

Control favors free driving for efficiency

# Experiments

## Compare **brake engage distance to endpoint**

control assuming delayed brakes

control assuming air brakes



## **Instantiate symbolic models** with concrete parameters from FRA

	<b>Parameter</b>	<b>Value</b>	<b>vs.</b>
$l_z$	Length	Medium = 2 345ft (40 cars)	short,long
$m$	Mass	Loaded = 10 520klb ( $263 \frac{\text{klb}}{\text{car}}$ )	empty
$v$	Speed	Fast = 60mph	slow
$F_{pb}$	Emergency brake	Loaded = 1 430klbf ( $35 750 \frac{\text{lb}}{\text{car}}$ )	empty,unknown
$t_{appl}$	Time	50 s	length-dependent
$A$	Acceleration	$5 \frac{\text{mph}}{\text{min}} = 391.91\text{klbf}$	
$f_a$	Brake force	$1.75 \frac{\text{mph}}{\text{min}} = 136.76\text{klbf}$	
$\epsilon$	Control cycle	100ms	

Brossaeu, J., Ede, B.M.: Development of an adaptive predictive braking enforcement algorithm. FRA/DOT/ORD-9/13 (2009)

# Experimental Results: Brake Engage Points

Cars	Slow						Fast					
	Loaded			Empty			Loaded			Empty		
	10	40	100	10	40	100	10	40	100	10	40	100
	<b>Brake force for unknown load <math>F_{pb} = 23\,338 \frac{\text{lbf}}{\text{car}}</math></b>											
Delay brakes	726	1,110	1,942	446	830	1,662	15,436	17,742	22,730	5,369	7,676	12,664
Air brakes	541	710	1,017	239	345	503	14,364	15,494	17,880	4,278	5,334	7,383
<b>Difference</b>	<b>185</b>	<b>400</b>	<b>925</b>	<b>207</b>	<b>485</b>	<b>1,161</b>	<b>1,072</b>	<b>2,248</b>	<b>4,850</b>	<b>1,091</b>	<b>2,342</b>	<b>5,281</b>
	<b>Brake force for known load, loaded: <math>F_{pb} = 35\,750 \frac{\text{lbf}}{\text{car}}</math>, empty: <math>F_{pb} = 10\,575 \frac{\text{lbf}}{\text{car}}</math></b>											
Delay brakes	597	982	1,814	554	939	1,771	10,817	13,123	18,111	9,277	11,583	16,571
Air brakes	409	565	822	364	512	746	9,743	10,859	13,188	8,200	9,309	11,602
<b>Difference</b>	<b>188</b>	<b>417</b>	<b>992</b>	<b>190</b>	<b>427</b>	<b>1,025</b>	<b>1,074</b>	<b>2,264</b>	<b>4,923</b>	<b>1,077</b>	<b>2,274</b>	<b>4,969</b>

# Experimental Results: Brake Engage Points



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	Loaded			Empty			Loaded			Empty		
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# Experimental Results: Brake Engage Points

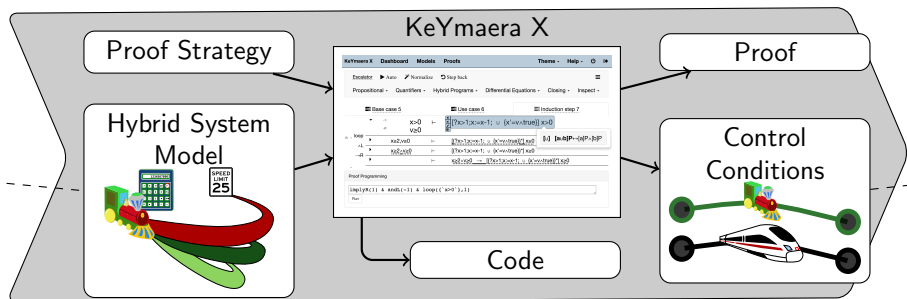
Air brake control conditions engage brakes considerably later

## Fast trains

	Cars	Loaded			Empty		
		10	40	100	10	40	100
Brake force for known load, loaded: $F_{pb} = 35\,750 \frac{\text{lbf}}{\text{car}}$ , empty: $F_{pb} = 10\,575 \frac{\text{lbf}}{\text{car}}$							
 Delay brakes		10,817	13,123	18,111	9,277	11,583	16,571
 Air brakes		9,743	10,859	13,188	8,200	9,309	11,602
<b>Difference</b>		<b>1,074</b>	<b>2,264</b>	<b>4,923</b>	<b>1,077</b>	<b>2,274</b>	<b>4,969</b>

**Difference exceeds FRA requirement of at most 1000ft undershoot!**

## Theorem proving ensures correct model



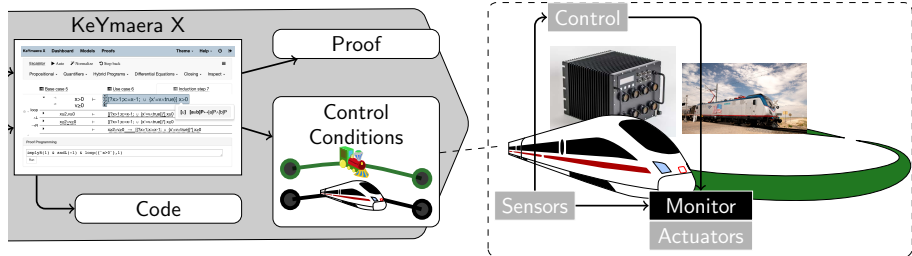
**No overshoot ✓ + Limited undershoot ✓**

**Proofs** for certification

**Models** and **code** for system architecture and implementation

**Control conditions** for runtime monitoring and testing

## Transfer safety of model to controller implementation



Monitor desired effect + safe environment

ModelPlex synthesizes and proves monitors for model compliance

Runtime: ensure safety and detect anomalies

Testing: generate and analyze test cases

Mitsch, S., Platzer, A.: ModelPlex: Verified runtime validation of verified cyber-physical system models. *Formal Methods in System Design*, 49(1), 2016.

KeYmaera X Dashboard Models Proofs Theme Help

Escalator Auto Normalize Step back

Propositional Quantifiers Hybrid Programs Differential Equations Closing Inspect

Base case 5 Use case 6 Induction step 7

	-1: $x > 0$	$\vdash$	$[[?x > 1; x := x - 1; \cup \{x' = v \wedge \text{true}\}]] x > 0$
	-2: $v \geq 0$		
loop	$x \geq 2, v \geq 0$	$\vdash$	$[[?x > 1; x := x - 1; \cup \{x' = v \wedge \text{true}\}]]^* x \geq 0$
$\wedge L$	$x \geq 2 \wedge v \geq 0$	$\vdash$	$[[?x > 1; x := x - 1; \cup \{x' = v \wedge \text{true}\}]]^* x \geq 0$
$\rightarrow R$		$\vdash$	$x \geq 2 \wedge v \geq 0 \rightarrow [[?x > 1; x := x - 1; \cup \{x' = v \wedge \text{true}\}]]^* x \geq 0$

Proof Programming

```
implyR(1) & andL(-1) & loop({'x>0'}, 1)
```

Run

[u] [a ub]P ↔ [a]P ∧ [b]P

[www.keymaeraX.org](http://www.keymaeraX.org)

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