

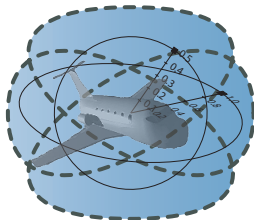
# 15-819/18-879: Logical Analysis of Hybrid Systems

## 03: Hybrid Systems

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- 1 Hybrid Systems
  - Labeled Transition Systems
  - Finite Automata
  - Hybrid Automata



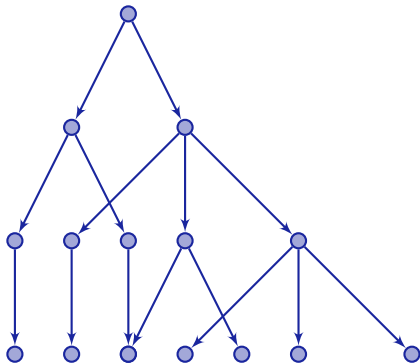
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## Definition (Labeled Transition System)

- Transition relation on  $Q \times A \times Q$ , denoted as  $q \xrightarrow{a} q^+$ , along with
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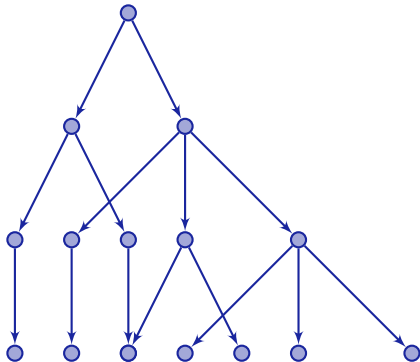
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## Definition (Trace)

(In)Finite series  $q_0, q_1, q_2, \dots \in Q$  such that for some  $a_i \in A$

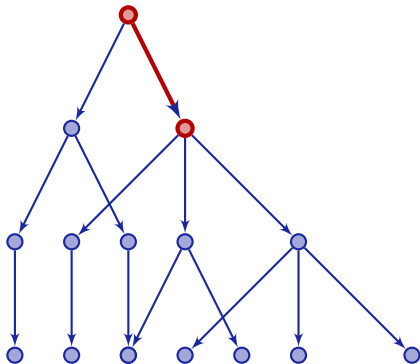
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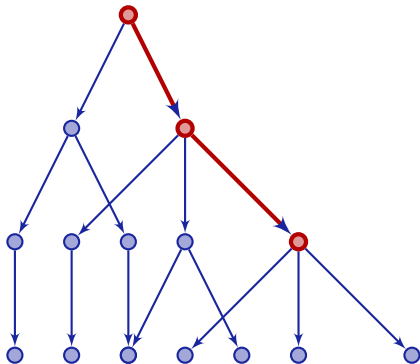
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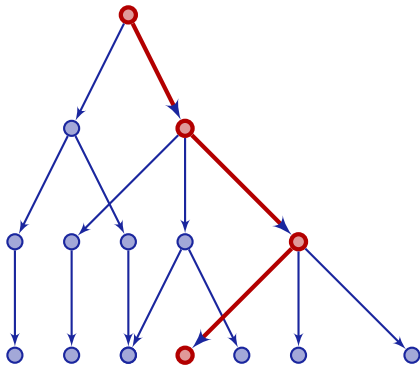




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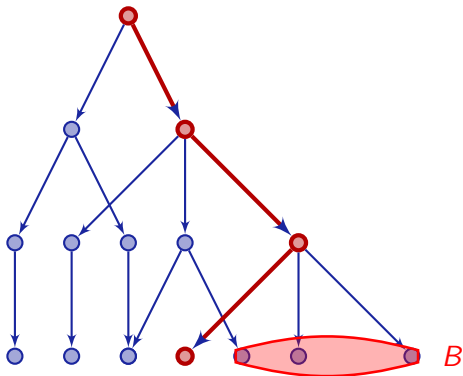
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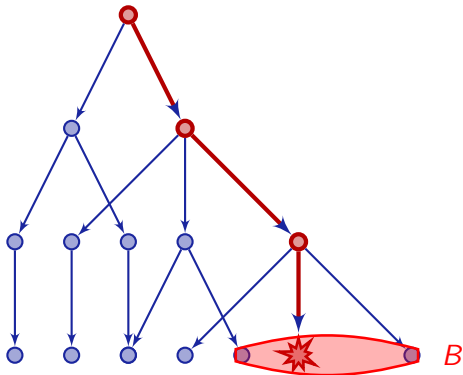
## Definition (Model Checking Problem for Reachability)

Given initial states  $Q_0 \subseteq Q$  and bad states  $B \subseteq Q$  for a transition system, check whether there is a trace from some  $q_0 \in Q_0$  to some  $q_b \in B$ .



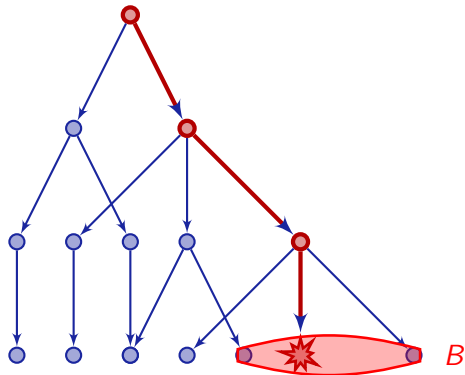
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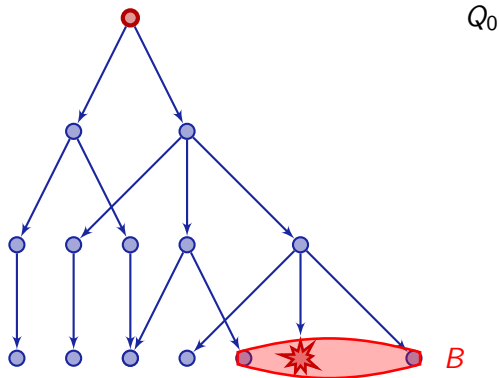
## Definition (Image Computation)

$$\text{Post}_A(Y) := \{q^+ \in Q : q \xrightarrow{a} q^+ \text{ for some } q \in Y, a \in A\}$$



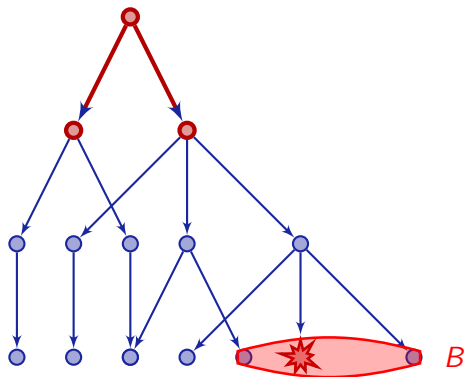
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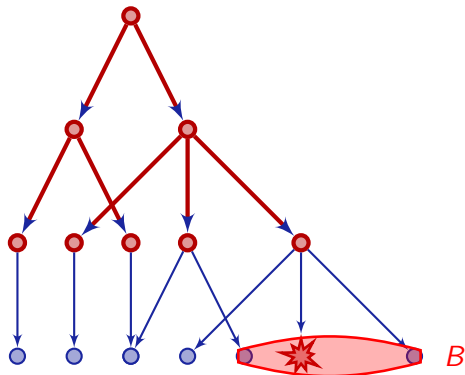
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$$Q_0 \xrightarrow{Post_A(Q_0)} Q_1 = Post_A(Q_0)$$

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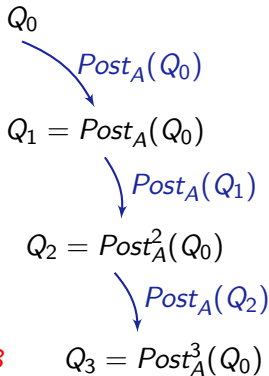
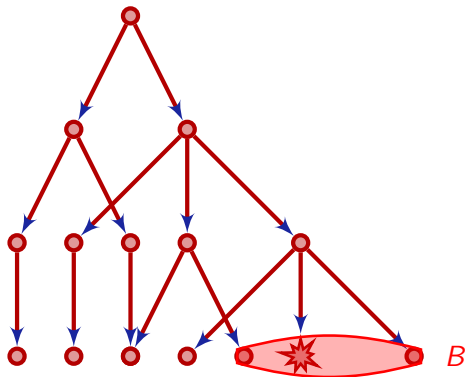
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$$\begin{array}{l}
 Q_0 \\
 \searrow^{Post_A(Q_0)} \\
 Q_1 = Post_A(Q_0) \\
 \searrow^{Post_A(Q_1)} \\
 Q_2 = Post_A^2(Q_0)
 \end{array}$$

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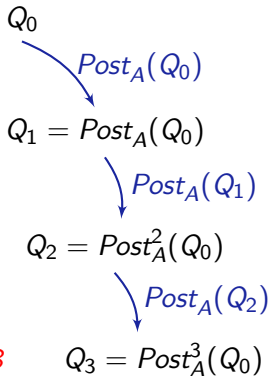
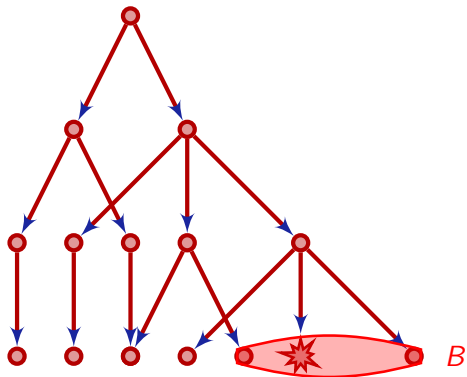




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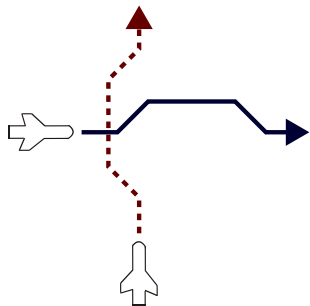
$$Post_A^*(Y) := \bigcup_{n \in \mathbb{N}} Post_A^n(Y) = \mu Z. (Y \cup Z \cup Post_A(Z))$$





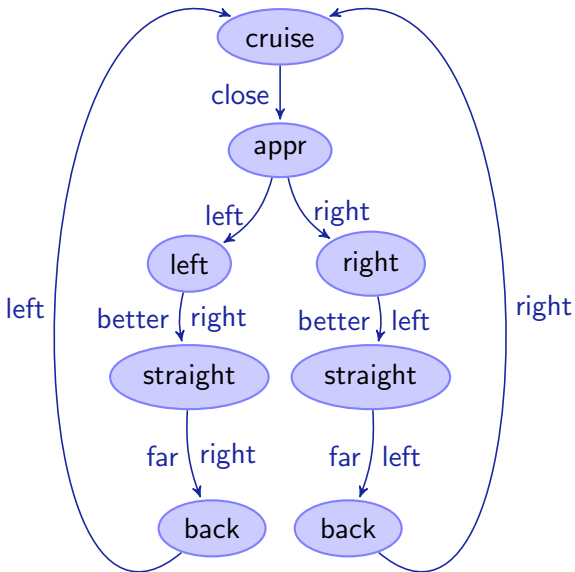
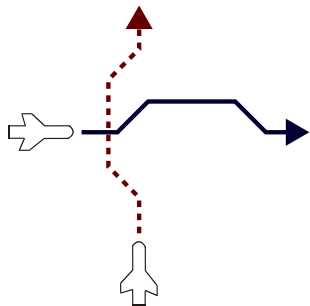
## Definition (Nondeterministic Finite Automata $\rightsquigarrow$ LTS)

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- finite set  $A$  of transition actions,
- finite set  $Q$  of states, initial states  $Q_0 \subseteq Q$ , accepting states  $F \subseteq Q$ .



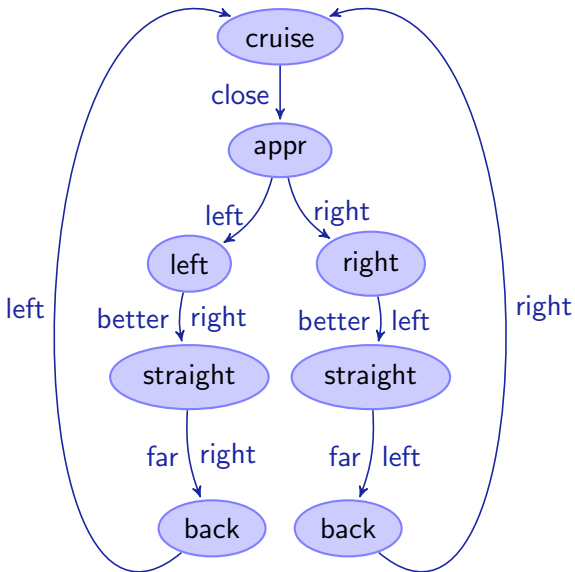
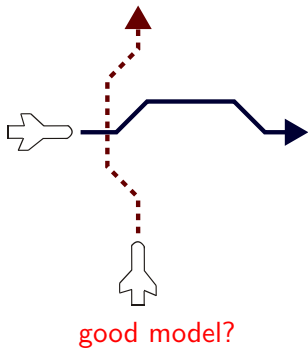


# Finite Automaton for Collision Avoidance





# Finite Automaton for Collision Avoidance



Collision avoidance is a property of  
controlled movement!

# Want Hybrid Transition Systems

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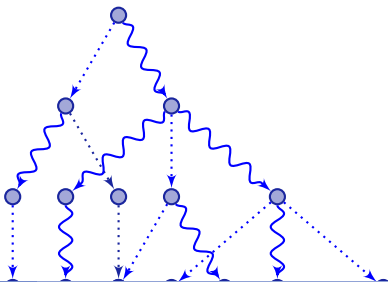
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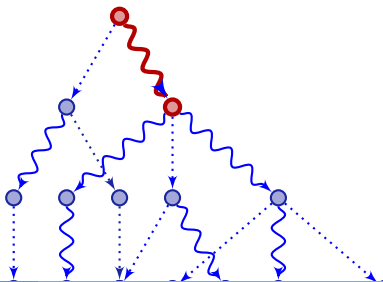


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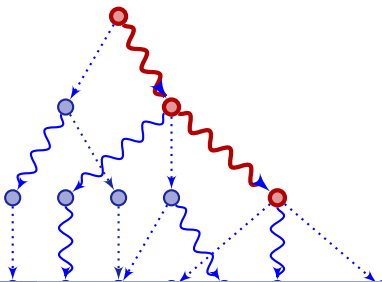


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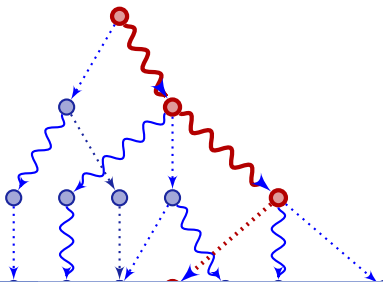


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Okay but how do we “justify” the respective transitions from our model??





“Graph switching between continuous dynamical systems”

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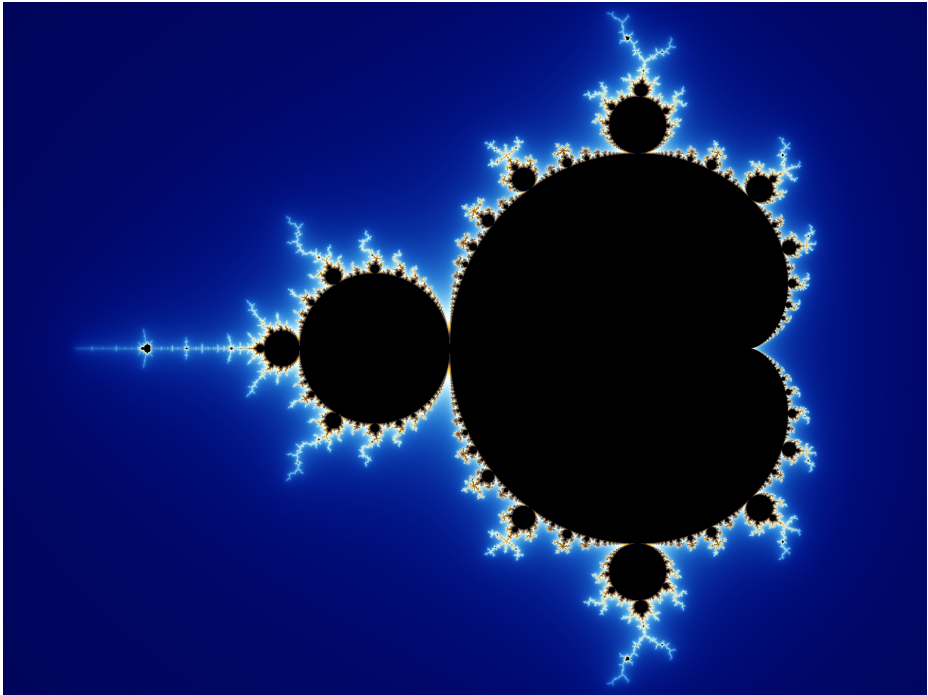
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usually comprising guard on current state and reset relations

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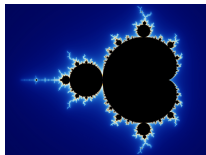
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Is this a good definition?



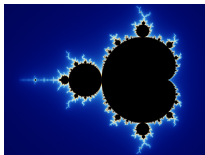




## Example (Mandelbrot Set)

For complex numbers  $c \in \mathbb{C}$  define  $f_0(c) = c$  and  $f_{n+1}(c) = f_n(c)^2 + c$ .  
Then the Mandelbrot set is

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Theorem (Lenore Blum, Cucker, Shub, Smale'90...98)

*“The Mandelbrot set is undecidable over  $\mathbb{R}$  / in Real Turing Machines”*

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What if  $inv_q$  is a Mandelbrot set?

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And all relations decidable / definable in first-order real arithmetic etc.

Computationally relevant output  
needs computational input!



M. S. Branicky, V. S. Borkar, and S. K. Mitter.

A unified framework for hybrid control: Model and optimal control theory.

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A. Platzer.

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Springer, Heidelberg, 2010.